

# MATH 285: Stochastic Processes

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## Today: MCMC

- Homework 3 is due on Friday, February 4, 11:59 PM

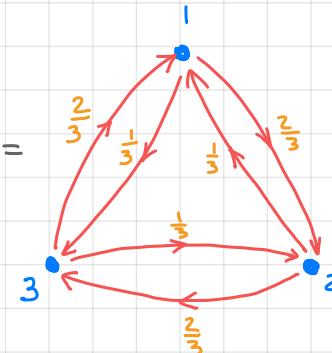
## Example

Consider random walk on  $G =$

Transition matrix

$$P = \begin{bmatrix} 0 & \frac{2}{3} & \frac{1}{3} \\ \frac{1}{3} & 0 & \frac{2}{3} \\ \frac{2}{3} & \frac{1}{3} & 0 \end{bmatrix},$$

$P$  is doubly stochastic, so  
 $\pi = \left( \frac{1}{3}, \frac{1}{3}, \frac{1}{3} \right)$



Detailed balance equation:

$$p(j,i) = \frac{\pi(j)}{\pi(i)} p(i,j) = p(i,j) \Rightarrow \text{not reversible}$$

If  $\pi = \left( \frac{1}{|S|}, \dots, \frac{1}{|S|} \right)$ ,  $(X_n)$  is reversible only if  $P = P^t$

## Example: Hard Core Configuration

Hard Core Configuration

on  $\{1, 2, \dots, N\}^2$  is a function

$$c: \{1, \dots, N\}^2 \rightarrow \{0, 1\}$$

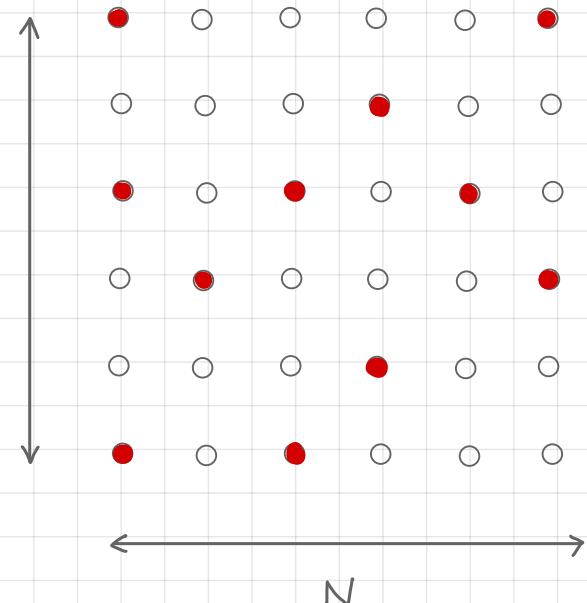
such that

$$c(i,j) = 1 \Rightarrow c(i \pm 1, j \pm 1) = 0$$

Denote by  $HCC_N$  the set of  
all hard core configurations

on  $\{1, \dots, N\}^2$ . Suppose we want to choose a uniform  
distribution on  $HCC_N$ ,  $P[Z = c] = \frac{1}{|HCC_N|} \quad \forall c \in HCC_N$

Problem: How to compute  $|HCC_N|$ ?



## Example: Hard Core Configuration

Computing  $|HCC_N|$  for large  $N$  is difficult.

Instead we construct a MC on  $HCC_N$  whose stationary distribution is the uniform distribution on  $HCC_N$ .

Construction: for any two configurations  $c \neq c'$

$$p(c, c') = \begin{cases} \frac{1}{N^2} & \text{if } c \text{ and } c' \text{ differ at exactly one point} \\ 0, & \text{otherwise} \end{cases}$$

$$p(c, c) = 1 - \sum_{c' \neq c} p(c, c')$$

Implementation: at each step choose  $(i, j) \in \{1, \dots, N\}^2$  uniformly at random and change the value at  $(i, j)$  if possible. E.g.,  $X_n = c$ , choose  $(i, j)$ .

- If  $c(i, j) = 1$ , then  $X_{n+1} = c'$  with  $c'(i, j) = 0$ ,  $c'(k, l) = c(k, l)$

## Example: Hard Core Configuration

- If  $c(i,j) = 0$ , and  $c(i \pm 1, j \pm 1) = 0$ , then  $X_{n+1} = c'$  with  
 $c'(i,j) = 1$  and  $c'(k,l) = c(k,l)$
- If  $c(i,j) = 0$  and one of  $c(i \pm 1, j \pm 1) \neq 0$ , then  $X_{n+1} = c$

- (i) Then for any  $c, c' \in HCC_N$   $\mathbb{P}[X_{n+1} = c' | X_n = c] = p(c, c')$
- (ii)  $(X_n)$  is irreducible ( $\forall c, c' \in HCC_N \quad p_{n_0}(c, 0) > 0, p_{n_1}(0, c') > 0$ )
- (iii)  $p(c, c') = p(c', c)$  ( $c$  and  $c'$  differ in only one coordinate)
- (iv) Uniform distribution on  $HCC_N$  is the stationary distribution
- $\pi(c) = 1/|HCC_N| \Rightarrow \pi(c) p(c, c') = \pi(c') p(c', c)$   
 $\Rightarrow \pi$  is stationary

Now if we start the process from any  $c \in HCC_N$ , then  
for sufficiently large  $n \quad \mathbb{P}[X_n = c] \approx \frac{1}{|HCC_N|}$

## Example : Graph coloring

Let  $G = (V, E)$  be a finite graph. A  $q$ -coloring of  $G$  (with  $q \in \mathbb{N}$ ) is a function  $f: V \rightarrow \{1, 2, \dots, q\}$  s.t.

$$u \sim v \Rightarrow f(u) \neq f(v)$$

(different colors of neighboring vertices)

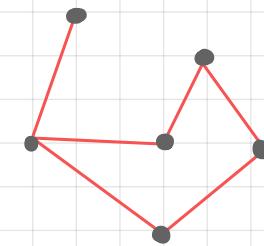
Q: How to choose a  $q$ -coloring uniformly at random?

Construct a MC: if  $f$  and  $g$  are two  $q$ -colorings,  $f \neq g$ ,

set  $p(f, g) = \begin{cases} \frac{1}{q|V|} & \text{if } f \text{ and } g \text{ differ at exactly one vertex} \\ 0 & \text{otherwise} \end{cases}$

$$p(f, f) = 1 - \sum_{g \neq f} p(f, g)$$

$(X_n)$  with transition probabilities  $p(f, g)$  is an irreducible MC with stationary distribution  $\pi(f) = \frac{1}{\#\{q\text{-coloring of } G\}}$



## Metropolis - Hastings Algorithm

Q: How to sample any (strictly positive) distribution  $\pi$ ?

Two-step MC: (1) propose moves (2) accept/reject move

### Construction of the Markov Chain

Let  $S$  be a finite set,  $\pi > 0$  a distribution on  $S$ .

(1) Construct an irreducible MC on  $S$  with symmetric transition probabilities  $q(i,j) = q(j,i)$ ,  $q(i,i) = 1 - \sum_{j \neq i} q(i,j)$

(2) If  $\pi$  (the desired distribution) is not uniform, construct a new MC with transition probabilities

$$p(i,j) = q(i,j) \min\left\{\frac{\pi(j)}{\pi(i)}, 1\right\}$$

- $\pi(i) p(i,j) = q(i,j) \min(\pi(j), \pi(i)) = q(j,i) \min\{\pi(i), \pi(j)\}$   
so  $\pi$  is stationary for  $p(i,j)$   $= \pi(j) p(j,i)$

## Metropolis - Hastings Algorithm

Suppose we know how to simulate a MC with transition probabilities  $q(i,j)$ . Then we can simulate a MC with transition probabilities  $p(i,j)$  using the two-step algorithm:

(i) Propose the move:

If  $X_n = i$ , then for  $j \neq i$  choose  $j$  with probability  $q(i,j)$

(ii) Accept or reject the move:

Accept the move with probability  $\min\left\{\frac{\pi(j)}{\pi(i)}, 1\right\}$

We get that  $P[X_{n+1} = j | X_n = i] = q(i,j) \min\left\{\frac{\pi(j)}{\pi(i)}, 1\right\} = p(i,j)$

If we now run  $(X_n)$  sufficiently long, then  $P[X_n = j] \approx \pi(j)$

Q: How long should we run  $(X_n)$ ?

## Convergence rate

Suppose that  $(X_n)$  is irreducible and aperiodic MC on  $S$  with  $|S|=N$ , and suppose that  $P$  is symmetric,  $P=P^t$ . Then by the spectral theorem  $P=UDU^t$

$$D = \begin{bmatrix} \lambda_1 & & \\ & \lambda_2 & \\ & \ddots & \\ & & \lambda_N \end{bmatrix}, \quad \lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_N$$

$$\begin{matrix} P^\infty \\ \| \end{matrix}$$

Perron - Frobenius theorem  $\Rightarrow \lim_{n \rightarrow \infty} P^n = U \lim_{n \rightarrow \infty} D^n U^t = U \begin{bmatrix} 1 & & \\ & \ddots & \\ & & 1 \end{bmatrix} U^t$

Then  $\|P^n - P^\infty\| = \max_{j \geq 2} |\lambda_j|^n$ . Mixing time:  $n$  s.t.  $\|P^n - P^\infty\|$  is small

E.g. for  $\begin{bmatrix} 1-\varepsilon & \varepsilon \\ \varepsilon & 1-\varepsilon \end{bmatrix} \quad \lambda_1 = 1, \quad \lambda_2 = 1-2\varepsilon \rightarrow (1-2\varepsilon)^n$  slow mixing

for  $\begin{bmatrix} \gamma_2 & \gamma_2 \\ \gamma_2 & \gamma_2 \end{bmatrix} \quad \lambda_1 = 1, \quad \lambda_2 = 0 \rightarrow \max_{j \geq 0} |\lambda_j|^n = 0$  fast mixing

## Example: Ising model

- $\Lambda_N = \{1, \dots, N\}^2$
- Spin configuration:  
 $\sigma : \Lambda_N \rightarrow \{-1, 1\}$
- Energy:  $H(\sigma) = - \sum_{\substack{i,j \\ i,j \in \Lambda}} \sigma(i)\sigma(j)$
- Gibbs measure:  $P_\beta(\sigma) = \frac{e^{-\beta H(\sigma)}}{Z_\beta}$
- where  $Z_\beta = \sum_{\sigma} e^{-\beta H(\sigma)}$  is the partition function (difficult)
- Take  $\pi(\sigma) = P_\beta(\sigma)$ . Then  $\pi(\sigma)/\pi(\sigma') = \exp(-\beta(H(\sigma) - H(\sigma')))$
- For  $\sigma \neq \sigma'$  take  $q(\sigma, \sigma') = \begin{cases} \frac{1}{N^2} & \text{if } \|\sigma - \sigma'\| = 2 \\ 0 & \text{otherwise} \end{cases}$
- Run MC ( $X_n$ ) with  $p(\sigma, \sigma') = q(\sigma, \sigma') \min\{1, \frac{\pi(\sigma')}{\pi(\sigma)}\}$

