

Let  $(X_t)_{t \in T}$  be a Markov process with (homogeneous) transition semigroup  $(q_t)_{t \in T}$ .  
If the initial distribution  $\text{Law}(X_0)$  is  $\mu$ , then for  $t > 0$ ,

$$\text{Law}(X_t)$$

**Def** A law  $\mu$  is called **invariant** or **stationary** if

(By the Markov property: if  $\mu$  is an invariant law and  $\text{Law}(X_{t_0}) = \mu$   
for some  $t_0 \in T$ , then  $\text{Law}(X_t) = \mu \quad \forall t \geq t_0$ .)

In the discrete-time setting, we only need to check the 1-step transition:

$$\text{Law}(X_{n+1})(dy) = \int \text{Law}(X_n)(dx) q(x, dy)$$

If the state space is discrete, this becomes a linear algebra question:

Let's focus on the finite state space, discrete-time setting,  $S = \{1, \dots, d\}$

$$\underline{v} = \begin{bmatrix} \mu(1) \\ \vdots \\ \mu(d) \end{bmatrix} \quad \mu(j) = \sum_i \mu(i) q(i, j) \quad P_{ij} = q(i, j)$$

Note: the row sums of  $P$  are  $1$ . Thus

$$P \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix} = \quad \text{so } (P - \mathbb{I}) \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix} = 0$$

We need a probability vector:  $v_j \geq 0$ ,  $\sum_j v_j = 1$ .

So long as  $w_j \geq 0$ , and  $\underline{w} \neq \underline{0}$  we can renormalize

Prop: Any finite state Markov chain has an invariant distribution.

Pf. We saw on the last slide that the transition matrix  $P$  possesses a vector  $\underline{w} \neq \underline{0}$  s.t.  $\underline{w}^T P = \underline{w}^T$ . Define  $\underline{v}$  by  $v_j :=$

Claim:  $\underline{v}^T P = \underline{v}^T$ .

$$\hookrightarrow v_j = |w_j| = \left| \sum_i w_i P_{ij} \right|$$

So: if  $\underline{v}^T P \neq \underline{v}^T$ , it follows that there is some  $j$  for which

$$\therefore \sum_j v_j$$

Thus, indeed,  $\underline{v}^T P = \underline{v}^T$ . Since  $\underline{v} \neq \underline{0}$  and  $v_j \geq 0 \forall j$ ,

$$\mu_j := v_j / \sum_i v_i$$

defines a probability vector, and

$$\mu^T P =$$

(This is a version of the Perron-Frobenius theorem.)

Note: the existence of an invariant distribution is not guaranteed in infinite state spaces.

Eg. Simple random walk on  $\mathbb{Z}$   $q(i, j) = \frac{1}{2} \mathbb{1}_{j=i+1} + \frac{1}{2} \mathbb{1}_{j=i-1}$   
If  $\mu$  satisfies  $\mu(j) = \sum_i \mu(i) q(i, j)$

$$\therefore \exists A, B \in \mathbb{R} \text{ s.t. } \mu(j) = A + Bj.$$

- Since  $\mu(j) \geq 0 \quad \forall j \in \mathbb{Z}$ ,
- Since  $\mu(j)$  must satisfy  $\sum_{j \in \mathbb{Z}} \mu(j) = 1$ ,

In fact, the random walk does have an

**invariant measure**  $\mu(j) = 1 \quad \forall j$

and it is unique up to positive multiple. But this

Markov chain does not have an invariant **probability** distribution.

Prop: Let  $(X_n)_{n \in \mathbb{N}}$  be a finite state, irreducible Markov chain.  
Then there is a unique invariant probability distribution  $\mu =$

$$\mu_i =$$

Pf. By first-step analysis

$$\mathbb{E}^i[\tau_j] = \sum_k q(i, k) \mathbb{E}^k[\tau_j(i, X)]$$

$$= 1 + \sum_{k \neq j} q(i, k) \mathbb{E}^k[\tau_j].$$

$\therefore$  if  $\mu$  is an invariant distribution, then  $\forall j$

$$\sum_i \mu_i \mathbb{E}^i[\tau_j] = \sum_i \mu_i + \sum_i \mu_i \sum_{k \neq j} q(i, k) \mathbb{E}^k[\tau_j]$$

we've shown that **if**  $\mu$  is an invariant probability distribution, then

$$\mu_i \cdot \mathbb{E}^i[\tau_i] = 1.$$

we already proved that an invariant distribution exists;  $\therefore$  this is it. **///**

**Cor:** In the finite state irreducible setting, the unique invariant distribution is strictly positive.

**Observation:** In [Lec. 44.1] we proved that  $0 < \mathbb{E}^i[\tau_i] < \infty \forall i$  (in the finite state irreducible case); we now have the additional fact:

A similar result holds for infinite chains: every irreducible, recurrent chain possesses a unique (up to scale) invariant measure, that is strictly positive.

Note: the equation  $E^i[\tau_j] = 1 + \sum_{k \neq j} q(i,k) E^k[\tau_j]$   
says that, if  $\underline{u}^{(j)}$  is the vector  $u_i^{(j)} =$  , then

$$\text{I.e. } \underline{u}^{(j)} = \begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} + P^{(j)} \underline{u}^{(j)}$$

This can be used to compute expected passage times,  
with linear algebra.

$$\begin{bmatrix} E^1[\tau_j] \\ \vdots \\ E^d[\tau_j] \end{bmatrix} = \underline{u}^{(j)}$$

See [Driver, § 22.9.2] for many worked examples.