

Markov chain:  $(X_n)_{n \in \mathbb{N}} : (\Omega, \mathcal{F}, \{\mathcal{F}_n\}_{n \in \mathbb{N}}, P) \rightarrow (S, \mathcal{B})$

Hitting time  $T_B(X) = \min\{n \geq 0 : X_n \in B\}$

We'd like to understand the distribution of  $T_B$ : e.g.  $P^x(T_B < \infty)$ ,  $E^x[T_B]$ .

We can study more general quantities;  
think in terms of games.

E.g.  $E^x[h(X_{T_B})]$

↳ start at  $x$ , continue until you hit  $B$ .

At the point  $y$  where you enter  $B$ ,  
collect \$  $h(y)$ .

E.g.  $E^x\left[\sum_{n=0}^{T_B} g(X_n)\right]$

↳ start at  $x$ , follow random path  $x_1, x_2, \dots, x_N$   
until  $x_N \in B$ . Pick up \$  $g(x_k)$  each time  
you visit  $x_k$ . How much do you collect total?

## Sub-Probability Kernels.

Let  $B \in \mathcal{B}$ . Recall that  $\mathcal{B}|_B := \{E \in \mathcal{B} : E \subseteq B\}$   
a  $\sigma$ -field over  $B$ .

If  $q: S \times \mathcal{B} \rightarrow [0, 1]$  is a probability kernel, we can restrict it to  $B$ :

$$q_B: B \times \mathcal{B}|_B \rightarrow [0, 1].$$

$$q_B(x, E) =$$

Note:  $\int q_B(x, dy) =$

As usual, we identify  $q_B$  with an operator

$$Q_B: \mathcal{B}(S, \mathcal{B}) \rightarrow \mathcal{B}(S, \mathcal{B})$$

$$(Q_B f)(x) = \int q_B(x, dy) f(y)$$

Has all the same regularity properties as  $Q$ ,  
except  $Q_B(1)$

Prop: Let  $h \in \mathcal{B}(S, \mathcal{B})$  or  $h: S \rightarrow [0, \infty]$   $\mathcal{B}/\mathcal{B}[0, \infty]$ -measurable.

Let  $(X_n)_{n \in \mathbb{N}}$  be a Markov process in  $(S, \mathcal{B})$  with 1-step transition operator  $Q$ . Let  $B \in \mathcal{B}$ . Define

$$u(x) := \mathbb{E}^x [h(X_{T_B}) : T_B < \infty], \quad x \in S.$$

Then  $u \in \mathcal{B}(S, \mathcal{B})$  or  $u: S \rightarrow [0, \infty]$  is  $\mathcal{B}/\mathcal{B}[0, \infty]$ -measurable,

and  $u = Qu$  on  $B^c$ ,  $u = h$  on  $B$ .

Moreover,  $u = Q_B h + Q_{B^c} u$  on  $B^c$ .

Pf. we use first step analysis:

$$\mathbb{E}^x [F(X_0, X_1, X_2, \dots)] = \int_S q(x, dy) \mathbb{E}^y [F(x, X_0, X_1, X_2, \dots)].$$

Here  $F(X_0, X_1, X_2, \dots) = h(X_{T_B}) \mathbb{1}_{T_B < \infty}$ .

Note, if  $x \in B^c$ , then  
and

$$\therefore F(x, X_0, X_1, X_2, \dots) =$$

$$\begin{aligned} \text{So } u(x) &= \mathbb{E}^x [F(X_0, X_1, X_2, \dots)] \\ &= \int_S q(x, dy) \mathbb{E}^y [F(x, X_0, X_1, X_2, \dots)] \end{aligned}$$

Prop: Let  $g: B^c \rightarrow [0, \infty]$  be  $\mathcal{B}/\mathcal{B}[0, \infty]$ -measurable. Set

$$v(x) := \mathbb{E}^x \left[ \sum_{0 \leq n < T_B} g(X_n) \right], \quad x \in S$$

Then  $v = 0$  on  $B$ ,  $v = Qv + g =$  on  $B^c$ .

Pf. Set  $F(X_0, X_1, X_2, \dots) = \sum_{0 \leq n < T_B} g(X_n)$ . Take  $x \in B^c$ .

$$\therefore F(x, X_0, X_1, X_2, \dots) =$$

$$\therefore v(x) = \mathbb{E}^x [F(X_0, X_1, X_2, \dots)] = \int_S q(x, dy) \mathbb{E}^y [F(x, X_0, X_1, X_2, \dots)]$$

So, we have to solve recursive equations

$$u = Q_{B^c} u + Q_B h \quad \text{on } B^c,$$

$$v = Q_{B^c} v + g \quad \text{on } B^c$$

If  $I - Q_B$  is invertible, and we can compute its inverse, we're good.

This is true if  $P^x(T_B < \infty) = 1 \quad \forall x \in B^c$  See [Driver, Cor 22.58]  
but otherwise is generally false.

In that case, we need to characterize which solution is the right one.

Theorem: In the problems considered in the last two propositions,  $u$  and  $v$  are the **unique minimal solutions** of the recursion equations.  
See [Driver, § 22.8]