

We've now seen that if $(Q_t)_{t \geq 0}$ is operator norm continuous @ $t=0$, then there is a generator A - a bounded operator on $B(S, \mathcal{B})$ s.t.

$$Q_t = e^{tA} = \sum_{n=0}^{\infty} \frac{t^n}{n!} A^n$$

One case where this level of regularity is commonplace is for discrete state spaces: S countable, $\mathcal{B} = 2^S$. Here

$$Q_t f(i) = \sum_{j \in S} q_t(i, j) f(j) \quad \forall i \in S$$

where $(q_t)_{t \geq 0}$ are the transition kernel mass functions

Note: $q_t(i, j) \in [0, 1]$, $\sum_{j \in S} q_t(i, j) = 1 \quad \forall i \in S$.

\therefore As a matrix, $\|q_t\|_{\infty} = \sup_{i \in S} \sum_{j \in S} q_t(i, j) = \sup_i 1 = 1$.

The continuity condition becomes

$$\begin{aligned} \|Q_t - I\|_{op} &= \|q_t - I\|_{\infty} = \sup_i \sum_j |q_t(i, j) - \delta_{ij}| \\ &= \sup_i (1 - q_t(i, i) + \sum_{j \neq i} q_t(i, j)) = \sup_i 2(1 - q_t(i, i)) \rightarrow 0 \text{ as } t \downarrow 0. \end{aligned}$$

Thus, if $(q_t)_{t \geq 0}$ are Markov transition kernel mass functions on a discrete state space S , satisfying

$$\lim_{t \downarrow 1} \inf_{i \in S} q_t(i, i) = 1$$

then the transition semigroup $(Q_t)_{t \geq 0}$ has a bounded generator A .

Question: what can we say about A ? $Q_t = e^{tA}$

$$\begin{aligned} Af(i) &= \left. \frac{d}{dt} \right|_{t=0^+} Q_t f(i) \\ &= \left. \frac{d}{dt} \right|_{t=0^+} \sum_{j \in S} q_t(i, j) f(j) \\ &\stackrel{?}{=} \sum_{j \in S} \left. \frac{d}{dt} \right|_{t=0^+} q_t(i, j) f(j) \end{aligned}$$

This suggests that A has a matrix a given by

$$a(i, j) = \left. \frac{d}{dt} q_t(i, j) \right|_{t=0^+}$$

If S is infinite, this takes some work to prove.

First question: does A even have a matrix?

You might think this must always be true: that every (bounded) linear operator

$$T: B(S) \rightarrow B(S)$$

has a matrix. Following the finite-dimensional case, we would expand

$$f = \sum_{j \in S} f(j) \mathbb{1}_{\{j\}}.$$

If f is a simple function, so $f(j) = 0$ for all but finitely many $j \in S$,

$$Tf(i) = \sum_{j \in S} f(j) T(\mathbb{1}_{\{j\}}(i))$$

So we would expect that T has matrix $\theta(i,j) = T(\mathbb{1}_{\{j\}})(i)$.

But if f is not simple, there's no way to extend this:

even if $\|\theta\|_\infty < \infty$, we can't check if $Tf(i) = \sum_j \theta(i,j)f(j)$.

Basic Problem: $f \equiv 1$ on \mathbb{N} .

$$f_k = \mathbb{1}_{\{1, 2, \dots, k\}}$$

$$\|f - f_k\|_\infty = 1.$$

Fact: \exists bounded operators on $B(S)$ that have no matrix.

Prop: Let $\{T_n\}_{n \in \mathbb{N}}$ be bounded operators on $B(S)$, each given by a matrix θ_n :

$$T_n f(i) = \sum_{j \in S} \theta_n(i, j) f(j). \quad \|\theta_n\|_\infty = \|T_n\|_{op} < \infty \quad \forall n.$$

If T is a bounded operator and $\|T_n - T\|_{op} \rightarrow 0$, then

T has a matrix θ given by $\theta(i, j) = \lim_{n \rightarrow \infty} \theta_n(i, j)$.

Pf. First, $T_n(\mathbb{1}_{\{j\}})(i) = \sum_{k \in S} \theta_n(i, k) \mathbb{1}_{\{j\}}(k) = \theta_n(i, j)$

$$\therefore \lim_{n \rightarrow \infty} \theta_n(i, j) = \lim_{n \rightarrow \infty} T_n(\mathbb{1}_{\{j\}})(i) = T(\mathbb{1}_{\{j\}})(i).$$

$$|T_n(\mathbb{1}_{\{j\}})(i) - T(\mathbb{1}_{\{j\}})(i)| \leq \sup_k |T_n(i, k) - T(i, k)| = \|T_n(\mathbb{1}_{\{j\}}) - T(\mathbb{1}_{\{j\}})\|_\infty \\ \leq \|T_n - T\|_{op} \|\mathbb{1}_{\{j\}}\|_\infty \rightarrow 0.$$

Now, $\{T_n\}_{n \in \mathbb{N}}$ is convergent, hence Cauchy.

$$(T_n - T_m) f(i) = T_n f(i) - T_m f(i) \\ = \sum_j \theta_n(i, j) f(j) - \sum_j \theta_m(i, j) f(j) \\ = \sum_j [\theta_n(i, j) - \theta_m(i, j)] f(j)$$

$\therefore T_n - T_m$ has matrix $\theta_n - \theta_m$, so

$$\|\theta_n - \theta_m\|_\infty = \|T_n - T_m\|_{op} \rightarrow 0 \text{ as } n, m \rightarrow \infty.$$

$$\sup_i \sum_j |\theta_n(i, j) - \theta_m(i, j)|$$

$$\therefore \text{For each } i, \sum_j |\theta_n(i, j) - \theta_m(i, j)| = \sum_j \liminf_{m \rightarrow \infty} |\theta_n(i, j) - \theta_m(i, j)|$$

$$\stackrel{\text{Fatou}}{\leq} \liminf_{m \rightarrow \infty} \sum_j |\theta_n(i, j) - \theta_m(i, j)| \leq \liminf_{m \rightarrow \infty} \|T_n - T_m\|_{op} = \|T_n - T\|_{op}.$$

As this is true for all i , taking $\sup_i \therefore \|\theta_n - \theta\|_\infty \leq \|T_n - T\|_{op}$

\therefore By the triangle inequality, $\|\theta\|_\infty < \infty$, so it defines a bounded operator

$$\hat{T}f(i) := \sum_j \theta(i, j) f(j).$$

$$\text{But then } \|\hat{T} - T_n\|_{op} = \|\theta - \theta_n\|_\infty \leq \|T - T_n\|_{op} \rightarrow 0$$

Since $\|T - T_n\|_{op} \rightarrow 0$, it follows that $T = \hat{T}$.

$$\text{I.e. } Tf(i) = \sum_j \theta(i, j) f(j).$$

///

Cor: Under the continuity condition $\|Q_t - I\|_{op} \rightarrow 0$ i.e. $\inf_i q_t(i,i) \rightarrow 1$ as $t \downarrow 0$, the semigroup has generator A with matrix a

$$a(i,j) = \left. \frac{d}{dt} \right|_{t=0^+} q_t(i,j).$$

Pf. We proved last lecture that $\|A - \frac{Q_t - I}{t}\|_{op} \rightarrow 0$ as $t \downarrow 0$.

Take any $t_n \downarrow 0$, $\therefore \|A - \frac{Q_{t_n} - I}{t_n}\|_{op} \rightarrow 0$

has matrix $\frac{1}{t_n} [q_{t_n}(i,j) - \delta_{ij}]$.

\therefore By our proposition, A has matrix

$$\begin{aligned} a(i,j) &= \lim_{n \rightarrow \infty} \frac{1}{t_n} [q_{t_n}(i,j) - \delta_{ij}] \\ &= \left. \frac{d}{dt} \right|_{t=0^+} q_t(i,j). \quad // \end{aligned}$$

Cor: $a(i,j) \geq 0$ for $i \neq j$, and $\sum_{j \in S} a(i,j) = 0 \quad \forall i \in S$.

(Exactly the same as the proof in [Lec.39.2].)