

Recurrence relations  
Math 475  
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Say we have a sequence of numbers  $a_0, a_1, a_2, \dots$  that satisfies a recurrence relation of the form

$$a_n = Ca_{n-1} + Da_{n-2}$$

whenever  $n \geq 2$  (here  $C, D$  are some constants and  $D \neq 0$ ). We want to find a closed formula for  $a_n$ .

The **characteristic polynomial** of this recurrence relation is defined to be

$$t^2 - Ct - D.$$

The roots of this polynomial are  $\frac{C \pm \sqrt{C^2 + 4D}}{2}$ . Call them  $\alpha$  and  $\beta$ . So we can factor the characteristic polynomial as

$$(1) \quad t^2 - Ct - D = (t - \alpha)(t - \beta).$$

Comparing constant terms, we get  $\alpha\beta = D$ , so  $\alpha \neq 0$  and  $\beta \neq 0$  because we assumed that  $D \neq 0$ .

Here is the first statement:

**Theorem 1.** *If  $\alpha \neq \beta$ , then there are constants  $c_0$  and  $c_1$  such that*

$$a_n = c_0\alpha^n + c_1\beta^n$$

for all  $n$ .

To solve for the coefficients, plug in  $n = 0$  and  $n = 1$  to get

$$\begin{aligned} a_0 &= c_0 + c_1 \\ a_1 &= \alpha c_0 + \beta c_1. \end{aligned}$$

Then you have to solve for  $c_0, c_1$  ( $a_0, a_1$  are part of the original sequence, so are given to you).

*Proof of Theorem 1.* Define a generating function

$$A(x) = \sum_{n=0}^{\infty} a_n x^n.$$

The recurrence relation says that we have a relation of the form

$$A(x) = CxA(x) + Dx^2A(x) + a_0 + (a_1 - Ca_0)x.$$

We can rewrite this as

$$(2) \quad A(x) = \frac{a_0 + (a_1 - Ca_0)x}{1 - Cx - Dx^2}.$$

We want to factor the denominator. To do this, plug in  $t \mapsto x^{-1}$  into (1) and multiply by  $x^2$  to get

$$1 - Cx - Dx^2 = (1 - \alpha x)(1 - \beta x).$$

Now we can apply partial fraction decomposition to (2) to write

$$A(x) = \frac{c_0}{1 - \alpha x} + \frac{c_1}{1 - \beta x}$$

for some constants  $c_0, c_1$ . But these terms are both geometric series, so we can further write

$$A(x) = c_0 \sum_{n=0}^{\infty} \alpha^n x^n + c_1 \sum_{n=0}^{\infty} \beta^n x^n.$$

The coefficient of  $x^n$  on the left side is  $a_n$  and the coefficient of  $x^n$  on the right side is  $c_0\alpha^n + c_1\beta^n$ . So we have equality for all  $n$ .  $\square$

There is a loose end: what if  $\alpha = \beta$ ?

**Theorem 2.** *If  $\alpha = \beta$ , then there are constants  $c_0$  and  $c_1$  such that*

$$a_n = c_0\alpha^n + c_1n\alpha^n$$

for all  $n$ .

Again, to solve for  $c_0, c_1$ , just plug in  $n = 0, 1$  to get a system of equations:

$$\begin{aligned} a_0 &= c_0 \\ a_1 &= c_0\alpha + c_1\alpha. \end{aligned}$$

*Proof.* We can start in the same way as in the previous proof. The only difference is that we are trying to take the partial fraction decomposition of

$$A(x) = \frac{a_0 + (a_1 - Ca_0)x}{(1 - \alpha x)^2}.$$

This can still be done, but now it looks like

$$\frac{d_0}{1 - \alpha x} + \frac{d_1}{(1 - \alpha x)^2}$$

for some constants  $d_0, d_1$ . The first is a geometric series, and the second we've seen: remember that  $1/(1 - x)^2 = \sum_{n \geq 0} (n + 1)x^n$ . So we get instead

$$A(x) = d_0 \sum_{n=0}^{\infty} \alpha^n x^n + d_1 \sum_{n=0}^{\infty} (n + 1)\alpha^n x^n.$$

Comparing coefficients, we get

$$a_n = d_0\alpha^n + d_1(n + 1)\alpha^n = (d_0 + d_1)\alpha^n + d_1n\alpha^n.$$

So  $c_0 = d_0 + d_1$  and  $c_1 = d_1$ .  $\square$

Let's finish with the example of the Fibonacci numbers  $f_n$ . These are defined by

$$\begin{aligned} f_0 &= 1 \\ f_1 &= 1 \\ f_n &= f_{n-1} + f_{n-2} \quad \text{for } n \geq 2 \end{aligned}$$

So the characteristic polynomial is  $t^2 - t - 1$ . Its roots are  $\frac{1 \pm \sqrt{5}}{2}$ . Set  $\alpha = (1 + \sqrt{5})/2$  and  $\beta = (1 - \sqrt{5})/2$ . So we have

$$f_n = c_0\alpha^n + c_1\beta^n$$

and we have to solve for  $c_0$  and  $c_1$ . Plug in  $n = 0, 1$  to get:

$$\begin{aligned} 1 &= c_0 + c_1 \\ 1 &= c_0\alpha + c_1\beta. \end{aligned}$$

So  $c_0 = 1 - c_1$ ; plug this into the second formula to get  $1 = (1 - c_1)\alpha + c_1\beta$ . Rewrite this as  $1 - \alpha = c_1(\beta - \alpha)$ . We can simplify this:  $\beta - \alpha = -\sqrt{5}$  and  $1 - \alpha = (1 - \sqrt{5})/2$ . So

$$c_1 = -\frac{1 - \sqrt{5}}{2\sqrt{5}}, \quad c_0 = 1 - c_1 = \frac{1 + \sqrt{5}}{2\sqrt{5}}.$$

In conclusion:

$$\begin{aligned} f_n &= \frac{1 + \sqrt{5}}{2\sqrt{5}} \left( \frac{1 + \sqrt{5}}{2} \right)^n - \frac{1 - \sqrt{5}}{2\sqrt{5}} \left( \frac{1 - \sqrt{5}}{2} \right)^n \\ &= \frac{1}{\sqrt{5}} \left( \frac{1 + \sqrt{5}}{2} \right)^{n+1} - \frac{1}{\sqrt{5}} \left( \frac{1 - \sqrt{5}}{2} \right)^{n+1}. \end{aligned}$$

(The last step wasn't necessary, we just did that to reduce the number of radical signs.)

What about higher degree recurrence relations like

$$a_n = C_1 a_{n-1} + C_2 a_{n-2} + \cdots + C_k a_{n-k} \quad \text{for } n \geq k?$$

This can be solved in the same way: one has to first find the roots of the characteristic polynomial  $t^k - C_1 t^{k-1} - C_2 t^{k-2} - \cdots - C_k$  and apply partial fraction decomposition. The simplest case is when the roots  $\alpha_1, \dots, \alpha_k$  are all distinct. In this case, we can say that there exist constants  $c_1, \dots, c_k$  such that

$$a_n = c_1 \alpha_1^n + \cdots + c_k \alpha_k^n$$

for all  $n$ . In order to solve for  $c_1, \dots, c_k$ , we have to consider  $n = 0, \dots, k - 1$  separately to get a system of  $k$  linear equations in  $k$  variables. When the roots appear with multiplicities, we have to do something like we did in Theorem 2. For example, if  $k = 5$  and the roots are  $\alpha$  with multiplicity 3 and  $\beta$  with multiplicity 2 (and  $\alpha \neq \beta$ ), then we would have

$$a_n = c_1 \alpha^n + c_2 n \alpha^n + c_3 n^2 \alpha^n + c_4 \beta^n + c_5 n \beta^n.$$

This should look familiar to you if you've ever solved a linear homogeneous differential equation with constant coefficients.

I'll leave it to you to formulate the general case, though we won't be doing anything more with it in this class.