# The Graph of Multiplication is Equivalent to Counting

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November 12, 1991

**Abstract**

Counting is $AC^0$-reducible to the graph of multiplication. Hence the graph of multiplication is equivalent under $AC^0$ reductions to majority and to the function form of multiplication.
*Keywords:* constant depth circuits, threshhold circuits, circuit complexity, graph of multiplication, computational complexity.

Chandra, Stockmeyer and Vishkin [4] showed that MULTIPLICATION and BINARY-COUNT and four other problems are equivalent with respect to constant-depth circuit reducibility (see also Furst, Saxe and Sipser [5]). In this note we prove that the graph of multiplication is also equivalent to these six problems, by proving that BINARY-COUNT is $AC^0$-reducible to the graph of multiplication.

The problems we are most interested in are:

MULTIPLICATION:
*Input*: Two integers $x$, $y$ in binary notation
*Output*: The product $x \cdot y$ in binary notation

MULTIPLICATION-GRAPH:
*Input*: Three integers $x$, $y$ and $z$ in binary notation
*Output*: Accept if and only if $x \cdot y = z$.

BINARY-COUNT:
*Input*: $x_1, \ldots, x_n \in \{0, 1\}$
*Output*: The binary representation of the number of $x_i$'s equal to 1

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The prior constant-depth circuit reduction of BINARY-COUNT to MULTIPLICATION is based on a method of Furst-Saxe-Sipser. For inputs $x_1, \ldots, x_n$, construct the integers $\overline{x}$, $\overline{y}$ with binary representations $\overline{x} = (x_1\vec{0}x_2\vec{0}\cdots x_{n-1}\vec{0}x_n)_2$ and $\overline{y} = (1\vec{0}1\vec{0}\cdots 1\vec{0}1)_2$ where $\vec{0}$ denotes $\log(n)$ many 0's. In the middle of the binary representation of the product $\overline{x}\cdot\overline{y}$ will be a block of $\log(n)$ many bits containing the binary representation of the sum of the $x_i$'s. However, the binary representation of $\overline{x}\cdot\overline{y}$ contains a lot of additional information; and it is easy to see that there is a constant depth circuit which, given $x_1, \ldots, x_n$ and given an integer $z$, determines if $z = \overline{x}\cdot\overline{y}$. Thus, the prior work left open the possibility that there are constant depth circuits for recognizing the graph of multiplication—this paper shows that such circuits do not exist.

Part of the reason we are interested in the non-existence of constant depth circuits for the graph of multiplication is that this implies the $\Delta_0^b$-predicates of bounded arithmetic [2] are not in $AC^0$; this should be contrasted with the work of Mancivi on $\Delta_0^b$-formulas in one free variable [7].

Recall the following theorem:

**Theorem 1** *(Chandra-Stockmeyer-Vishkin [4])*

(a) *MULTIPLICATION* $\leq_{AC^0}$ *BINARY-COUNT*

(b) *BINARY-COUNT* $\leq_{AC^0}$ *MULTIPLICATION*

Obviously MULTIPLICATION-GRAPH $\leq_{AC^0}$ MULTIPLICATION. For the converse reduction we shall prove:

**Theorem 2** *BINARY-COUNT* $\leq_{AC^0}$ *MULTIPLICATION-GRAPH.*

Recall that $TC^0$ is the set of predicates which can be computed by a family of unbounded fanin, constant depth circuits formed from NOT gates, AND gates and MAJORITY gates. In [4] it is also shown that MAJORITY is $AC^0$-equivalent to MULTIPLICATION, etc.

**Corollary 3** *The decision problem MULTIPLICATION-GRAPH is* $AC^0$ *equivalent to MULTIPLICATION and to BINARY-COUNT. Unbounded fanin, constant depth circuits with* NOT *gates,* AND *gates and gates for the graph of multiplication capture exactly* $TC^0$.

**Proof** The proof of Theorem 2 involves a modification of the Furst-Saxe-Sipser method. Now we let

$$\overline{x} = (x_1 0^{k_1-1} x_2 0^{k_2-1} x_3 \cdots 0^{k_{n-1}-1} x_n)_2$$

and

$$\overline{y} = (10^{k_{n-1}-1}10^{k_{n-2}-1}1 \cdots 0^{k_1-1}1)_2$$

where $0^k$ denotes $k$ occurences of the symbol 0. As before the product $\overline{x} \cdot \overline{y}$ will contain a block of bits encoding the sum of the $x_i$'s; however, the values of $k_i$ will be chosen that the rest of the product is easily predictable and thus can not aid in determining the sum of the $x_i$'s.

Let $z$ be the (binary representation of the) product $\overline{x} \cdot \overline{y}$. The computation of $z$ by the usual multiplication algorithm reduces to computing the summation of the values $\overline{x} \cdot 2^{m_j}$ where $m_j = \sum_{i=1}^{j} k_i$ for $0 \leq j < n$. We shall choose values for $k_i$ so that each column of this summation has at most a single "1" except, of course, that the middle column will contain all the bits $x_1, \ldots, x_n$. In addition, each $k_j$ will be greater than $\log n$ and thus the sum of the $x_i$'s in the middle column will not propogate any carries into a column containing a "1". Let $C = \sum_{i=0}^{n-1} k_i$, it follows that columns $C$ through $C + \log n$ of $z$ contain the binary representation of the number of $x_i$'s equal to one. (Columns are numbered from right-to-left starting with 0.)

We define

$$k_j \;=\; n^2 + j.$$

Since $x_j$ occurs in column $\sum_{i=j}^{n-1} k_i$ of $\overline{x}$ and since $\overline{y}$ has a "1" in precisely columns $\sum_{i=1}^{m-1} k_i$ for $1 \leq m \leq n$, the summation for $z$ contains the bit $x_j$ in the columns numbered

$$g(j,m) = \sum_{i=1}^{m-1} k_i + \sum_{i=j}^{n-1} k_i$$

for $1 \leq m \leq n$. When $j = m$, $g(j,m) = C$; so each $x_j$ occurs in column $C$ of the summation. To show that in all other columns there is at most a single non-zero bit, it suffices to show that if $(j,m)$ and $(j',m')$ are distinct pairs with $j \neq m$ then $g(j,m) \neq g(j',m')$. Without loss of generality, $m' < m \leq j < j'$ and thus it suffices to show that

$$\sum_{i=m'}^{m-1} k_i \;\neq\; \sum_{i=j}^{j'-1} k_i$$

i.e., that

$$\sum_{i=m'}^{m-1} n^2 + i \;\neq\; \sum_{i=j}^{j'-1} n^2 + i.$$

But since $\sum_{i=1}^{n-1} i < n^2$, it must be that the two summations have the same number of terms; which is impossible.

We have shown that for $\overline{x}$ and $\overline{y}$ as above, the product $z$ contains the number of nonzero $x_i$'s in columns $C$ through $C + \log n$; the other bits of $z$ are zero except for that for all $j \neq m$, $z$ has bit value $x_j$ in column $g(j, m)$. Thus BINARY-COUNT can be reduced to MULTIPLICATION-GRAPH by in parallel trying every possible number of nonzero values of $x_i$'s and using graph-of-multiplication gates to determine if the numbers are correct. □

Let $LH$ denote the logtime hierarchy, i.e., the set of predicates accepted by logtime, constant alternation Turing machines. $LH$ was originally defined by Sipser and is a uniform version of $AC^0$; see [1, 3] for more information. To conclude, we observe that the constant depth circuits reducing BINARY-COUNT to the graph of multiplication are $LH$-uniform. To see this, it suffices to note that since

$$g(j, m) = n^2(n + k - j) + \frac{k^2 - k}{2} + \frac{n^2 - n}{2} - \frac{j^2 - j}{2}$$

$g(j, m)$ can be computed by constant depth, polynomial size, uniform circuits. This is because the multiplication function is in constant alternation, linear time (see Lipton [6] or Wilkie [8]) and because the computation of $g(j, m)$ involves addition and multiplication of numbers of length $O(\log n)$.

**Acknowledgements**

This construction is based, in part, on an idea due to P. Pudlák. The final form was constructed during conversations with E. Allender and H. Straubing.

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