

**Math 160B - Winter 2022 - Homework 3**  
**Due Wednesday, January 26, 11:00pm**  
(Hand in by uploading to Gradescope)

You can also find these problems in the PDF text.

1. Let  $R$  be a binary relation on  $\Sigma^*$ . Define  $S$  to be the set

$$S := \{w \in \Sigma^* : \exists v \in \Sigma^* (\langle v, w \rangle \in R)\}.$$

We call  $S$  the *projection* of  $R$  onto its second coordinate. Suppose that  $R$  is c.e. Prove that  $S$  is c.e.

2. Suppose  $R$  and  $S$  are  $k$ -relations on  $\Sigma^*$ . Define the *symmetric difference*  $R \Delta S$  of  $R$  and  $S$  to equal  $(R \setminus S) \cup (S \setminus R)$ . Suppose that  $R \Delta S$  is finite. Prove that  $R$  is decidable if and only if  $S$  is decidable. Prove that  $R$  is c.e. if and only if  $S$  is c.e.

3. Suppose that  $R \subseteq \mathbb{N}$  is enumerated by an algorithm  $M$  in increasing order. In other words,  $M$  successively outputs  $n_1, n_2, \dots$  so that  $R = \{n_1, n_2, n_3, \dots\}$  and  $n_i < n_{i+1}$  for each  $i$ . Prove that  $R$  is decidable. What if the assumption is weakened to the condition that  $n_i \leq n_{i+1}$ . Must  $R$  still be decidable?

4. Prove that a  $k$ -ary relation is computably enumerable if and only if it is the domain of a partial computable function.

5. Prove that a set is computably enumerable if and only if it is empty or is the range of a recursive function.

6. Prove Craig's Theorem:

- (a) Suppose  $\Gamma$  is c.e. set of propositional formulas. Prove there is a decidable set  $\Pi$  of propositional formulas such that  $\Gamma \models \Pi$ .
- (b) Suppose  $\Gamma$  is c.e. set of first-order sentences. Prove there is a decidable set  $\Pi$  of first-order sentences such that  $\Gamma \models \Pi$ .

Hint: The proofs for (a) and (b) are identical. One method of proof is to enumerate  $\Gamma$  as  $A_1, A_2, \dots$  and consider the formulas  $A_1 \wedge A_2 \wedge \dots \wedge A_k$ .