

**Start Time:**

**Your name:**

**Stop Time:**

**Integrity signature:**

Time limit 15 minutes, not counting download and upload. Please add explanation if over 17 minutes.

**There are two pages, the second with “cheat” sheet material. Only the first page needs to be turned in.**

1. Prove that there exists an FO proof of

$$\forall x (P(x) \rightarrow Q(x)) \rightarrow \exists x P(x) \rightarrow \exists x Q(x).$$

## Axioms and Inference Rules for FO-proofs

**PL1:**  $A \rightarrow B \rightarrow A$

**PL2:**  $(A \rightarrow B \rightarrow C) \rightarrow (A \rightarrow B) \rightarrow (A \rightarrow C)$

**PL3:**  $\neg A \rightarrow A \rightarrow B$

**PL4:**  $(\neg A \rightarrow A) \rightarrow A$

**EQ1:**  $x = x$

**EQ2:**  $x = y \rightarrow y = x$

**EQ3:**  $x = y \rightarrow y = z \rightarrow x = z$

**EQ<sub>f</sub>:**  $y_1 = z_1 \rightarrow \dots \rightarrow y_k = z_k \rightarrow f(y_1, \dots, y_k) = f(z_1, \dots, z_k)$

**EQ<sub>P</sub>:**  $y_1 = z_1 \rightarrow \dots \rightarrow y_k = z_k \rightarrow P(y_1, \dots, y_k) \rightarrow P(z_1, \dots, z_k)$

**Universal Instantiation (UI):**  $\forall x A(x) \rightarrow A(t)$ .

**Modus Ponens:**  $\frac{A \quad A \rightarrow B}{B}$ .

**Generalization (Gen):**  $\frac{C \rightarrow A}{C \rightarrow \forall x A}$  ( $x$  not free in  $C$ )

$A \vee B$ ,  $A \wedge B$ , and  $\exists x A$  abbreviate  $\neg A \rightarrow B$ ,  $\neg(A \rightarrow \neg B)$  and  $\neg \forall x \neg A$ .

## Constructing FO-proofs

**Derived rules:**

**Propositional rules:** Modus Tollens, Hypothetical Syllogism and more generally Tautological Implication (TAUT).

**Generalization:**  $\frac{A}{\forall x A}$

If  $t$  is substitutable for  $x$  in  $A$ :

**UI Rule:**  $\frac{\forall x A(x)}{A(t)}$       **Substitution:**  $\frac{A(x)}{A(t)}$

For Deduction/Contradiction/By Cases, assume  $A$  is a sentence.

**Deduction Theorem:**  $\Gamma \vdash A \rightarrow B$  iff  $\Gamma, A \vdash B$ .

**Proof by Contradiction:**

$\Gamma \vdash A$  iff  $\Gamma \cup \{\neg A\}$  is inconsistent.

$\Gamma \vdash \neg A$  iff  $\Gamma \cup \{A\}$  is inconsistent.

**Proof by Cases:** If  $\Gamma, A \vdash B$  and  $\Gamma, \neg A \vdash B$ , then  $\Gamma \vdash B$ .

**Theorem on Constants.**  $c$  is a new constant symbol, not in  $\Gamma, A, B$ .

For  $A(x)$  a formula,  $\Gamma \vdash \forall x A(x)$  iff  $\Gamma \vdash A(c)$ .

If  $\exists x A(x)$  is a sentence,  $\Gamma \cup \{\exists x A(x)\} \vdash B$  iff  $\Gamma \cup \{A(c)\} \vdash B$