Success and challenges in determining the rational points on curves



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Diophantine equations

Example problems: Find the solutions $x, y \in \mathbb{Q}$ to

$$\begin{aligned} x^2 + y^2 &= 1\\ x^2 + y^2 &= -1\\ x^2 + y^2 &= 5\\ x^2 + y^2 &= 3\\ 3x^3 + 4y^3 &= 5\\ x^6 + 8x^5 + 22x^4 + 22x^3 + 5x^2 + 6x + 1 &= y^2\\ x^6 + x^2 + 1 &= y^2\\ x^6 + 6x^5 - 15x^4 + 20x^3 + 15x^2 + 30x - 17 &= y^2\\ (x^3 - x^2 - 2x + 1)y^7 - (x^3 - 2x^2 - x + 1) &= 0\\ x^4 + y^4 + x^2y + 2xy - y^2 + 1 &= 0\\ x^2y^2 - xy^3 - x^3 - 2x^2 + y^2 - x + y &= 0 \end{aligned}$$

Note: All of these ask for the *rational points* on curves.

Definition: A curve *C* over \mathbb{Q} is *nice* if it is:

smooth, projective, absolutely irreducible. **Typical example:** Smooth plane projective curve:

$$C: X^4 + Y^4 + X^2 YZ + 2XYZ^2 - Y^2Z^2 + Z^4 = 0$$

Decision problem: Given a nice curve C over \mathbb{Q} ,

decide if $C(\mathbb{Q}) = \emptyset$.

Determination problem: Given a nice curve C over \mathbb{Q} ,

find a useful description of $C(\mathbb{Q})$.

For curves of genus > 1: List the finite set $C(\mathbb{Q})$.

- 1. Outline of a procedure to tackle the decision problem
- 2. Highlight challenges in executing the procedure
- 3. Finite Descent as a tool to face these challenges
- 4. Results for smooth plane quartics

Adelic points:

$$C(\mathbb{Q}) \hookrightarrow C(\mathbb{A}) := C(\mathbb{R}) \times \prod_p C(\mathbb{Q}_p)$$

Global-Local principle:

 $C(\mathbb{Q}) \neq \emptyset$ implies $C(\mathbb{A}) \neq \emptyset$

Happy fact: Deciding if $C(\mathbb{A}) = \emptyset$ is decidable.

Local-Global principle fails:

 $C(\mathbb{A}) \neq \emptyset$ does not imply $C(\mathbb{Q}) \neq \emptyset$,

Examples:

$$3X^3 + 4Y^3 + 5Z^3 = 0$$

$$X^4 + Y^4 + X^2 YZ + 2XYZ^2 - Y^2Z^2 + Z^4 = 0$$

Alternative approach: Embed curve C in another variety with a sparser set of rational points, e.g., an Abelian variety J.

Theorem (Mordell-Weil): $J(\mathbb{Q})$ is a finitely generated abelian group:

$$J(\mathbb{Q}) \simeq \underbrace{J(\mathbb{Q})_{\mathrm{tors}}}_{\mathrm{finite}} \times \mathbb{Z}^r$$

Principal homogeneous space: $C \subset \underline{\operatorname{Pic}}_{C}^{1}$ under $J = \underline{\operatorname{Pic}}_{C}^{0}$.

 $\underline{\operatorname{Pic}}_{C}^{1}(\mathbb{Q}) \neq \emptyset$ if and only if $\underline{\operatorname{Pic}}_{C}^{1} \simeq J$

Challenge: Decide if $\underline{\operatorname{Pic}}_{C}^{1}(\mathbb{Q}) = \emptyset$ or find $\mathfrak{d} \in \underline{\operatorname{Pic}}_{C}^{1}(\mathbb{Q})$.

If $\underline{\operatorname{Pic}}_{C}^{1}(\mathbb{Q}) = \emptyset$ then $C(\mathbb{Q}) = \emptyset$. Otherwise $\iota_{\mathfrak{d}} \colon C \hookrightarrow J$.

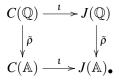
Challenge: Compute $J(\mathbb{Q}) \simeq J(\mathbb{Q})_{\text{tors}} \times \mathbb{Z}^r$, in particular *r*.

Mordell-Weil group combined with adelic information

Assume:

- We have $\mathfrak{d} \in \underline{\operatorname{Pic}}^1_C(\mathbb{Q})$.
- We have generators for $J(\mathbb{Q})$.

Commutative diagram:



(Watch the Poonen • which modifies the $J(\mathbb{R})$ factor)

Conjecture: Writing $\overline{C(\mathbb{Q})} \subset C(\mathbb{A})$ for the topological closure,

 $\overline{C(\mathbb{Q})} \stackrel{?}{=} \iota(C(\mathbb{A})) \cap \overline{\widetilde{
ho}(J(\mathbb{Q}))}$

(see [Scharaschkin, B-Elkies (ANTS V), Flynn, B.-Stoll])

$$C(\mathbb{Q}) \xrightarrow{\iota} J(\mathbb{Q})/BJ(\mathbb{Q})$$

$$\downarrow^{\rho_{S}} \qquad \qquad \downarrow^{\rho_{S}}$$

$$\prod_{p \in S} C(\mathbb{F}_{p}) \xrightarrow{\iota_{S}} \prod_{p \in S} J(\mathbb{F}_{p})/B \cdot \operatorname{im}(\rho_{p})$$

Let S be a finite set of primes ; B a positive integer

► Let
$$\Lambda_p = \ker(\rho_p : J(\mathbb{Q}) \to J(\mathbb{F}_p))$$
 and $\Lambda_S := \bigcap_{p \in S} \Lambda_p$
► $C(\mathbb{Q}) \to V_{S,B} := \operatorname{im}(\iota_S) \cap \operatorname{im}(\rho_S) \subset \frac{J(\mathbb{Q})}{\Lambda_S + BJ(\mathbb{Q})}$

Heuristic (Poonen): For appropriate *S*, *B*, the set $V_{S,B}$ consists only of cosets containing a point from $C(\mathbb{Q})$.

INPUT: A nice curve *C* of genus g > 0. **OUTPUT:** $P \in C(\mathbb{Q})$ or Unsolvable if $C(\mathbb{Q}) = \emptyset$. Execute in parallel:

0. Try candidates for $P \in C(\mathbb{Q})$ and return P if one is found. Information from $V_{S,B}$ (step 5) helps.

and

- 1. If $C(\mathbb{A}) = \emptyset$ return Unsolvable
- 2. Determine $\mathfrak{d} \in \underline{\operatorname{Pic}}^1_{\mathcal{C}}(\mathbb{Q})$ or return Unsolvable if $\underline{\operatorname{Pic}}^1_{\mathcal{C}}(\mathbb{Q}) = \emptyset$.
- 3. Determine $J(\mathbb{Q})$.
- 4. Choose reasonable values for *S*,*B*.
- 5. Mordell-Weil sieving: If $V_{S,B} = \emptyset$ return Unsolvable.
- 6. Increase S, B; go to 5.

Test case (B.-Stoll): Consider genus 2 curves admitting a model

$$C: y^2 = f_6 x^6 + f_5 x^5 + \dots + f_0 \text{ with } f_i \in \{-3, \dots, 3\}$$

Success: We were able to decide for all of them!

All curves	196171	100.00%
Curves with rational points	137 490	70.09%
Curves without rational points	58681	29.91 %
Curves with $C(\mathbb{A}) \neq \emptyset$	166768	85.01 %
Curves with $C(\mathbb{A}) \neq \emptyset$ and $C(\mathbb{Q}) = \emptyset$	29278	14.92%
Curves that need BSD conjecture	42	0.02%

Disclosure: We only really needed MW-sieving for 1445 of these curves (27786 of these curves have a non-trivial 2-cover obstruction to having rational points)

How to deal with rational points

(see [Chabauty, Coleman, Flynn]) **Problem:** If $P \in C(\mathbb{Q})$ then $V_{S,B}$ is never empty.

Idea (Chabauty): Construct a *p*-adic analytic function Φ_p on $C(\mathbb{Q}_p)$ that vanishes on $C(\mathbb{Q})$.

Restriction: Construction only works if $rkJ(\mathbb{Q}) = r < g$.

Sketch of procedure:

1. Use MW-Sieving to find S, B and $P_i \in C(\mathbb{Q})$ such that

$$V_{S,B} = \{P_1,\ldots,P_n\} + \Lambda_S + BJ(\mathbb{Q})$$

2. Find prime p with $BJ(\mathbb{Q}) \subset \Lambda_p$ such that

 $P_i \not\equiv P_j \pmod{p}$ for any $i \neq j$

3. For each P_i , use Φ_p to show that there are no other rational points Q with $Q \equiv P_i \pmod{p}$

Computational Challenges

No guarantee that either procedure will terminate, i.e.:

- We only have a heuristic that MW-sieving converges to a sharp result.
- We have no guarantee we can always find a *p* such that Φ_p does not have inconvenient extraneous *p*-adic zeros.

Bigger problem: we cannot guarantee we can get started:

For decision procedure:

- Decide if $\underline{\operatorname{Pic}}_{C}^{1}(\mathbb{Q}) = \emptyset$ or find $\mathfrak{d} \in \underline{\operatorname{Pic}}_{C}^{1}(\mathbb{Q})$.
- Determine the *r* in $J(\mathbb{Q}) \simeq J(\mathbb{Q})_{\text{tors}} \times \mathbb{Z}^r$
- ► Find generators for J(Q)

For determination procedure:

• What to do if $r \ge g$?

(See [Wetherell, B.; future: Kim, Balakrishnan?])

Multiplication-by-n:

$$0 \to J[n] \to J \xrightarrow{n} J \to 0$$

Taking galois cohomology:

$$0 \to \frac{J(\mathbb{Q})}{nJ(\mathbb{Q})} \xrightarrow{\gamma} H^1(\mathbb{Q}, J[n]) \to H^1(\mathbb{Q}, J)$$

Approximate image locally:

Explicit descent computations: We need to work with

$$\gamma \colon \frac{J(k)}{nJ(k)} \to H^1(k, J[n]) \text{ for } k = \mathbb{Q}, \mathbb{R}, \mathbb{Q}_p$$

- ► How do we represent *J*(*k*)?
- How do we represent $H^1(k, J[n])$?
- How do we compute γ?

Representing J(k):

 $\operatorname{Pic}^{0}(C/k) \subset J(k)$; equality if $C(\mathbb{A}) \neq \emptyset$. Use divisors on the curve.

Problem: We only know how to efficiently represent $H^1(k, M)$ for a very limited class of Galois modules.

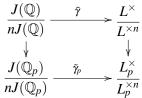
Twisted power: Let *M* be a Galois module and $\Delta = \operatorname{Spec} L = \{\theta_1, \dots, \theta_m\}$ a Galois set. Define

 $M^{\Delta}:=M\theta_1\oplus\cdots\oplus M\theta_m$

Hilbert 90: $H^1(k, \mu_n^{\Delta}) = L^{\times}/L^{\times n}$. Let $J[n] = \operatorname{Spec}(L)$. Consider $0 \to J[n] \to (\mu_n)^{J[n]} \to R^{\vee} \to 0$ Cohomology: $H^1(k, J[n]) \to L^{\times}/L^{\times n}$.

Computations using descent setups

(see [Cassels, Schaefer, Poonen-Shaefer, B.-Poonen-Stoll]) Writing $L_p = L \otimes \mathbb{Q}_p$



- Map $\tilde{\gamma}$ is induced by a function $f \in k(C) \otimes L$.
- Images of $\tilde{\gamma}_p$ are computable.
- ► For most *p*, this image lands in "unramified" part
- Image of $\tilde{\gamma}$ is generated by *S*-units.

$$\operatorname{Sel}^{\tilde{\gamma}}(J) = \{ \delta \in L^{\times}/L^{\times n} : \rho_p(\delta) \in \operatorname{im} \tilde{\gamma}_p \text{ for all } p \}$$

Bounding Ranks:

$$\frac{J(\mathbb{Q})}{nJ(\mathbb{Q})} = \frac{J(\mathbb{Q})_{\text{tors}}}{nJ(\mathbb{Q})_{\text{tors}}} \times \left(\frac{\mathbb{Z}}{n\mathbb{Z}}\right)'$$

So bounding the size of $im \gamma$ bounds r (hopefully sharply).

Embedding curve in *J*:

$$[\underline{\operatorname{Pic}}_{C}^{1}] \in H^{1}(\mathbb{Q}, J[2g-2])$$

There exists $\mathfrak{d} \in \underline{\operatorname{Pic}}^1_C(\mathbb{Q})$ if and only if $[\underline{\operatorname{Pic}}^1_C] \in \operatorname{im} \gamma$.

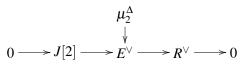
Bonus: Map $\tilde{\gamma}$ can be evaluated immediately on *C*.

$$\mathrm{Sel}^{\tilde{\gamma}}(C) = \{ \delta \in L^{\times}/L^{\times n} : \rho_p(\delta) \in \tilde{\gamma_p}(C(\mathbb{Q}_p)) \text{ for all } p \}$$

Example: Smooth plane quartics (B.-Poonen-Stoll)

Let *C* be a smooth plane quartic.

- Set $\Delta = \operatorname{Spec}(L)$ of 28 bitangents
- Even weight vectors $E \subset (\mathbb{Z}/2\mathbb{Z})^{\Delta}$:



Cohomology:

$$\begin{split} \frac{J(k)}{2J(k)} & \xrightarrow{\tilde{\gamma}} \frac{L^{\times}}{L^{\times 2}k^{\times}} \\ & \downarrow^{\gamma} & \downarrow^{\gamma} \\ 0 & \rightarrow J[2](k) & \rightarrow E^{\vee}(k) & \rightarrow R^{\vee}(k) & \rightarrow H^{1}(J[2]) & \rightarrow H^{1}(E^{\vee}) \end{split}$$

We need all of computational algebraic number theory ...

$$\begin{split} \frac{J(k)}{2J(k)} &\xrightarrow{\tilde{\gamma}} \frac{L^{\times}}{L^{\times 2}k^{\times}} \\ & \downarrow^{\gamma} & \bigvee^{\gamma} \\ 0 & \rightarrow J[2](k) & \rightarrow E^{\vee}(k) & \rightarrow R^{\vee}(k) & \rightarrow H^{1}(J[2]) & \rightarrow H^{1}(E^{\vee}) \end{split}$$

- $\tilde{\gamma}$ consists of evaluation at the "generic" bitangent.
- ▶ We need the ring of integers of *L* and *S*-units in *L*.
- ► $J[2](k), R^{\vee}(k), E^{\vee}(k)$ follow from identifying

 $\operatorname{Gal}(L/k) \subset \operatorname{Sp}_6(\mathbb{F}_2).$

Theorem: Consider

$$C: X^{3}Y - X^{2}Y^{2} - X^{2}Z^{2} - XY^{2}Z + XZ^{3} + Y^{3}Z = 0.$$

Then $J(\mathbb{Q}) \simeq \mathbb{Z}/51\mathbb{Z}$ and

 $C(\mathbb{Q}) = \{(1:1:1), (0:1:0), (0:0:1), (1:0:0), (1:1:0), (1:0:1)\}.$

Theorem: Consider

$$C: X^2Y^2 - XY^3 - X^3Z - 2X^2Z^2 + Y^2Z^2 - XZ^3 + YZ^3 = 0.$$

Assuming GRH, we have $J(\mathbb{Q}) \simeq \mathbb{Z}$ and

$$\begin{split} C(\mathbb{Q}) &= \{(1:1:0), (-1:0:1), (0:-1:1), (0:1:0), \\ &\quad (1:1:-1), (0:0:1), (1:0:0), (1:4:-3)\}. \end{split}$$

Observation: The map $\tilde{\gamma}$ can be evaluated on *C* directly.

$$ilde{\gamma} \colon C(\mathbb{Q}) o rac{L^{ imes}}{L^{ imes 2} \mathbb{Q}^{ imes}}$$

Comparing local images gives another computable obstruction to rational points.

Theorem: Consider

$$C: X^4 + Y^4 + X^2 YZ + 2XYZ^2 - Y^2Z^2 + Z^4 = 0$$

Then $C(\mathbb{A}) \neq \emptyset$ but assuming GRH one can prove that *C* has no rational points.

Kiran, Everett, Joe, Organizing committee, Program committee THANK YOU!!

For a wonderful

