

On the decomposition of graphs into complete bipartite subgraphs

by

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Abstract

For a given graph G , we consider a \mathbf{B} -decomposition of G , i.e., a decomposition of G into complete bipartite subgraphs G_1, \dots, G_t , such that any edge of G is in exactly one of the G_i 's. Let $\alpha(G; \mathbf{B})$ denote the minimum value of $\sum_i |V(G_i)|$ over all \mathbf{B} -decompositions of G . Let $\alpha(n; \mathbf{B})$ denote the maximum value of $\alpha(G; \mathbf{B})$ over all graphs on n vertices.

A \mathbf{B} -covering of G is a collection of complete bipartite subgraphs G'_1, G'_2, \dots, G'_t , such that any edge of G is in at least one of the G'_i . Let $\beta(G; \mathbf{B})$ denote the minimum value of $\sum_i |V(G'_i)|$ over all \mathbf{B} -coverings of G and

let $\beta(n; \mathbf{B})$ denote the maximum value of $\beta(G; \mathbf{B})$ over all graphs on n vertices.

In this paper, we show that for any positive ε , we have

$$(1 - \varepsilon) \frac{n^2}{2e \log n} < \beta(n; \mathbf{B}) \leq \alpha(n; \mathbf{B}) < (1 + \varepsilon) \frac{n^2}{2 \log n}$$

where $e = 2.718\dots$ is the base of natural logarithms, provided n is sufficiently large.

Introduction

For a finite graph G , a *decomposition* P of G is a family of subgraphs G_1, G_2, \dots, G_t , such that any edge in G is an edge of exactly one of the G_i 's. If all G_i 's belong to a specified class of graphs \mathbf{H} , such a decomposition will be called an \mathbf{H} -decomposition of G (see [2]).

Let f denote a *cost function* for graphs which assigns certain non-negative real values to all graphs. Sometimes it is desirable to decompose a given graph into subgraphs in \mathbf{H} such that the total "cost" (the sum of the cost function values of all subgraphs) is minimized. In other words, for a given graph G , we consider the following:

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$$\alpha_f(G; \mathbf{H}) = \min_P \sum_i f(G_i)$$

where $P = \{G_1, G_2, \dots, G_i\}$ ranges over all \mathbf{H} -decompositions of G .

Also of interest to us will be the quantity

$$\alpha_f(n; \mathbf{H}) = \max_G \alpha_f(G; \mathbf{H})$$

where G ranges over all graphs on n vertices.

If we take f_0 to be the counting function, which assigns value 1 to any graph, and \mathbf{P} is the family of all planar graphs, then $\alpha_{f_0}(G; \mathbf{P})$ is simply the thickness of G . If \mathbf{F} denotes the family of forests, then $\alpha_{f_0}(G; \mathbf{F})$ is called the arboricity of G (see [6]). Many results along these lines are available. The reader is referred to [2] for a brief survey.

In this paper, we will deal almost exclusively with the case in which \mathbf{H} is \mathbf{B} , the family of complete bipartite graphs. By a theorem in [5], the value of $\alpha_{f_0}(n; \mathbf{B})$ is given by:

$$\alpha_{f_0}(n; \mathbf{B}) = n - 1.$$

We consider the cost function f_1 where the value $f_1(G)$ is just the number of vertices in G . In the remaining part of the paper, we abbreviate $\alpha(n) = \alpha_{f_1}(n; \mathbf{B})$ and $\alpha(G) = \alpha_{f_1}(G; \mathbf{B})$. In particular, we show for any given ε and sufficiently large n ,

$$(1) \quad (1 - \varepsilon) \frac{n^2}{2e \log n} < \alpha(n) < (1 + \varepsilon) \frac{n^2}{2 \log n}$$

where e satisfies $\ln e = 1$.

An \mathbf{H} -covering of G is a collection of subgraphs of G , say G'_1, \dots, G'_i , such that any edge of G is in at least one of the G'_i , and all G'_i are in \mathbf{H} . For a given cost function f , we can define

$$\beta_f(G; \mathbf{H}) = \min_P \sum_i f(G'_i)$$

where $P = \{G'_1, \dots, G'_i\}$ ranges over all \mathbf{H} -coverings of G .

It is easily seen that

$$\beta_f(G; \mathbf{H}) \leq \alpha_f(G; \mathbf{H})$$

and

$$\beta_f(n; \mathbf{H}) \leq \alpha_f(n; \mathbf{H}).$$

We will show that the asymptotic growth of $\beta_{f_1}(n; \mathbf{B})$ is quite similar to $\alpha_{f_1}(n; \mathbf{B})$. In fact, we will obtain the same upper and lower bounds for $\beta_{f_1}(n; \mathbf{B})$ as those for $\alpha_{f_1}(n; \mathbf{B})$ in (1).

A lower bound

We derive these bounds mainly by probabilistic methods, which have been extensively described in the book by two of the authors [4].

Theorem 1. $\alpha(n) \geq (1 - \varepsilon) \frac{n^2}{2e \log n}$ for any given positive ε and sufficiently large n .

Proof. Let us consider a random graph G with n vertices and $\lfloor n^2/2e \rfloor$ edges. The probability of G containing a complete bipartite subgraph $K_{a,b}$ is bounded above by

$$\binom{n}{a} \binom{n}{b} e^{-ab} < e^{(a+b) \log n - ab}$$

(where $\lfloor x \rfloor$ and $\lceil x \rceil$ denote the greatest integer less than x and the least integer greater than x , respectively.)

Let S denote the set of all unordered pairs $\{a, b\}$ satisfying

$$1 \leq a, b \leq n, \frac{a+b}{ab} < \frac{1-\varepsilon}{\log n}.$$

The probability of G containing one of the complete bipartite subgraphs $K_{a,b}$ with $\frac{a+b}{ab} < \frac{1-\varepsilon}{\log n}$ is bounded above by

$$\sum_{\{a,b\} \in S} \binom{n}{a} \binom{n}{b} e^{-ab} < \sum_{\{a,b\} \in S} e^{-\varepsilon ab} < \sum_{\{a,b\} \in S} e^{-\varepsilon (\log n)^2} < n^2 e^{-\varepsilon (\log n)^2} < 1$$

for large n .

Therefore, there exists a graph G with n vertices and $\lfloor n^2/2e \rfloor$ edges such that G does not contain any $K_{a,b}$ as a subgraph. Let $P = \{G_1, G_2, \dots, G_t\}$ denote a \mathbf{B} -decomposition of G such that $\alpha(G)$ is the sum of the sizes of vertex set $V(G_i)$ of G_i , i.e.,

$$\alpha(G) = \sum_{i=1}^t |V(G_i)|.$$

For any edge (u, v) in G , we define

$$f(u, v) = \frac{|V(G_i)|}{|E(G_i)|}$$

where $\{u, v\}$ is in $E(G_i)$, the edge set of G_i .

It is easily seen that

$$\alpha(G) = \sum_{\{u,v\}} f(u, v).$$

Since G does not contain $K_{a,b}$ as a subgraph, any $G_i = K_{c,d}$, $1 \leq i \leq t$, satisfies that $\frac{c+d}{cd} \geq \frac{1-\varepsilon}{\log n}$. Thus we have

$$f(u, v) \geq \frac{1-\varepsilon}{\log n} \text{ for any } \{u, v\} \text{ in } E(G).$$

and

$$\alpha(n) > \alpha(G) > \frac{(1-\varepsilon)n^2}{2e \log n}$$

for sufficiently large n . This proves the theorem.

An upper bound

First, we shall prove a preliminary result.

Lemma. For any $\varepsilon > 0$ any graph on n vertices and $\rho \binom{n}{2}$ edges contains a complete bipartite graph $K_{s,t}$ as a subgraph where $t = \lfloor 1(1-\varepsilon)n\rho^s \rfloor$ and $s < \varepsilon\rho n$ for n sufficiently large.

Proof. Suppose G has n vertices and $\rho \binom{n}{2}$ edges and G does not contain $K_{s,t}$ as a subgraph. From the proof in [3], the following holds:

$$(2) \quad n(\rho n - s)^s \leq (t-1) \cdot n^s.$$

However, on the other hand, we have

$$(t-1)n^s < tn^s \leq (1-\varepsilon)n^{1+s}\rho^s < n(\rho n - s)^s$$

since $s < \varepsilon\rho n$.

This contradicts (2). Thus G must contain $K_{s,t}$.

Theorem 2. For any given ε , we have

$$(3) \quad \alpha(n) < (1+\varepsilon) \frac{n^2}{2 \log n}$$

if n is large enough.

Proof. From Lemma 1, one can easily verify that a graph G on $\rho \binom{n}{2}$ edges and n vertices contains a subgraph H isomorphic to $K_{s,t}$, where $s = \lfloor (1-\varepsilon_1) \log n \log(1/\rho) \rfloor$ and $t = \lfloor s^2 \log(1/\rho) \rfloor$ and $\varepsilon_1 > \frac{(\log n)^2}{\rho n}$. We will decompose G into complete bipartite subgraphs by a "greedy algorithm". Given G we find a subgraph H isomorphic to $K_{s,t}$ and let G_1 to be the subgraph of G containing all edges of G except those in H . Now, we find a subgraph H_1 isomorphic to K_{s_1,t_1} and let G_2 to be a subgraph of G_1 containing all

edges of G_1 except those in H_1 and continue in this fashion until only $\varepsilon_2 \frac{n^2}{\log n}$ edges are left. Thus G is decomposed into H, H_1, \dots , together with $\varepsilon_2 \frac{n^2}{\log n}$ edges and we have the following recursive relation

$$(4) \quad \alpha(G) \leq s + t + \alpha(G_1).$$

We will prove by induction that for a give $\varepsilon < \varepsilon_2 < \varepsilon_1, \varepsilon_3 > 0$ and sufficiently large n the following holds,

$$(5) \quad \alpha(G) \leq (1 + \varepsilon_2) \frac{n^2}{2 \log n} \int_0^\rho \log(1/x) dx + 2\varepsilon_2 \frac{n^2}{\log n}.$$

Suppose (5) holds for any graph H with $|E(H)| < \rho \binom{n}{2}$. From (4), we have

$$\alpha(G) \leq (1 - \varepsilon_2) (\log n)^2 / (\log(1/\rho))^3 + (1 + \varepsilon_2) \frac{n^2}{2 \log n} \int_0^{\rho'} \log(1/x) dx + 2\varepsilon_2 \frac{n^2}{\log n}$$

where $\rho' = (|E(G)| - st) / \binom{n}{2}$ for n sufficiently large. It suffices to show that

$$\begin{aligned} (1 - \varepsilon_2) (\log n)^2 / \log(1/\rho)^3 + (1 + \varepsilon_2) \frac{n^2}{2 \log n} \int_0^{\rho'} \log(1/x) dx &\leq \\ &\leq (1 + \varepsilon_2) \frac{n^2}{2 \log n} \int_0^{\rho'} \log(1/x) dx \end{aligned}$$

This can be verified by straightforward calculation. Thus (5) is proved and we have

$$\alpha(n) \leq (1 + \varepsilon_2) \frac{n^2}{2 \log n} \int_0^1 \log(1/x) dx + 2\varepsilon_2 \frac{n^2}{\log n} \leq (1 + \varepsilon) \frac{n^2}{2 \log n}$$

for given $\varepsilon > 0$. Theorem 2 is proved.

By slightly modifying the proofs of Theorem 1, we can easily prove the following.

Theorem 3.

$$\beta_{f_1}(n; \mathbf{B}) \geq (1 - \varepsilon) \frac{n^2}{2e \log n}$$

for any positive ε and sufficiently large n .

Therefore we have

$$(1 - \varepsilon) \frac{n^2}{2e \log n} < \beta_{f_t}(n; \mathbf{B}) \leq \alpha_{f_t}(n; \mathbf{B}) < (1 + \varepsilon) \frac{n^2}{2 \log n}$$

for any given positive ε and sufficiently large n , which summarizes the main results of the paper.

Some related question

As we noted earlier, the lower bound is obtained by a probabilistic method which is nonconstructive. It would be of great interest to find an explicit construction of a graph G on n vertices, $c_1 n^2 / \log n$ edges (or $c_2 n^2$ edges) which does not contain an $K_{c_3 \log n, c_3 \log n}$ as a subgraph for some constants c_1, c_2 and c_3 .

Another interesting problem which has long been conjectured [4] concerns the Turán number $T(K_{t,t}; n)$, the maximum number of edges a graph on n vertices can have which does not contain $K_{t,t}$ as a subgraph. Is it true that

$$T(K_{t,t}; n) = O(n^{2-1/t})?$$

For the case $t = 3$, the above equality has been verified in [1].

In this paper, we have shown that $\alpha_{f_t}(n; \mathbf{B}) = O(n^2 / \log n)$. However, we do not know the existence of

$$\lim_{n \rightarrow \infty} \frac{\alpha_{f_t}(n; \mathbf{B})}{n^2 / \log n} \quad \text{or} \quad \lim_{n \rightarrow \infty} \frac{\beta_{f_t}(n; \mathbf{B})}{n^2 / \log n},$$

obviously.

Let \mathbf{G}_n be the set of all the $2^{\binom{n}{2}}$ labelled graphs on n vertices. It would be of interest to evaluate $\sum_{G \in \mathbf{G}_n} \alpha_{f_t}(G; \mathbf{B})$. It is not unreasonable to conjecture that

$$\lim_{n \rightarrow \infty} \frac{\sum_{G \in \mathbf{G}_n} \alpha_{f_t}(G; \mathbf{B})}{2^{\binom{n}{2}} n^2 / \log n} = c$$

exists and c is probably equal to $\lim_{n \rightarrow \infty} \frac{\alpha_{f_t}(n; \mathbf{B})}{n^2 / \log n}$. We can also ask the analogous question for $\beta_{f_t}(G; \mathbf{B})$.

Let $\mathbf{G}_{n,m}$ be the set of all graphs on n vertices and m edges. We can define $\alpha_{f_t}(n, m; \mathbf{H})$ to be the maximum value of $\alpha_{f_t}(G; \mathbf{H})$ where G ranges over all graphs in $\mathbf{G}_{n,m}$. In this paper we investigate $\alpha_{f_t}(n, m; \mathbf{B})$ where m is about $n^2 / 2e$. One could also investigate $\alpha_{f_t}(n, m; \mathbf{B})$ or $\beta_{f_t}(n, m; \mathbf{B})$. In particular, we can ask the problem of determining m so that $\alpha(n, m; \mathbf{B})$ is maximized or to find the range for m for which we have $\alpha(n, m; \mathbf{B}) = o(n^2)$.

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