

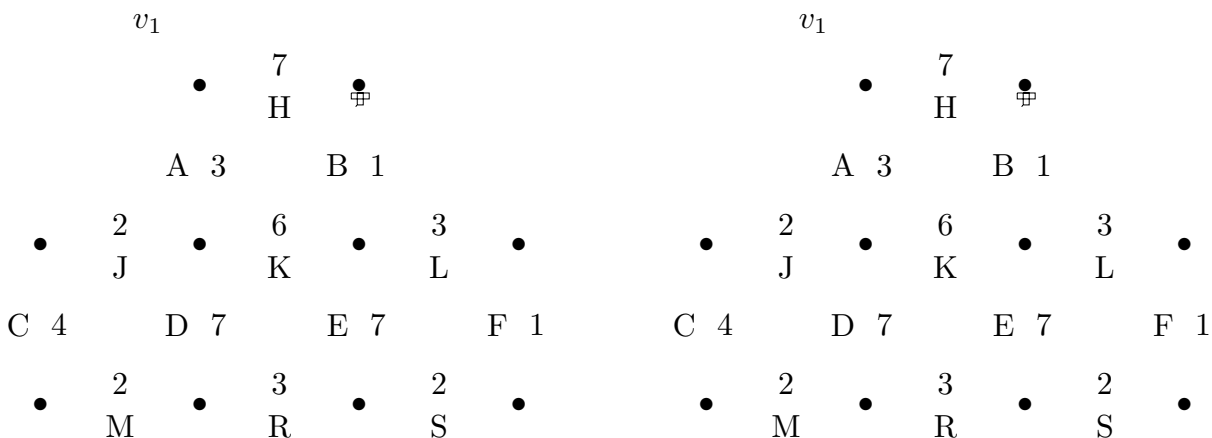
Name \_\_\_\_\_ ID No. \_\_\_\_\_

**IF YOU WANT YOUR GRADE WITHOUT THE FINAL:  
CROSS OUT PROBLEM 1, HAND IN YOUR EXAM, AND TELL US.**

There are 210 points possible.

1. (30 pts.) Recall that Dijkstra’s algorithm finds shortest paths from  $v_1$  to all other vertices by adding edges linking in the closest points. In the graph shown below, each edge is bidirectional; that is, you can travel in either direction on it. **Edges are labeled with upper case letters.** (Two copies of the graph are provided so you can use one as a “worksheet” if you wish.)

- (a) List edges in order chosen by algorithm: \_\_\_\_\_
- (b) At each vertex, give the length of the shortest path from  $v_1$  to the vertex. *Indicate which graph has your answer.*



2. (25 pts.) Consider the following eight complexity categories (remember  $\lg = \log_2$ ):

$$\Theta(n) \quad \Theta(n^2) \quad \Theta(2^n) \quad \Theta(3^{\lg n}) \quad \Theta(n^{\lg n}) \quad \Theta(n \lg n) \quad \Theta((\sqrt{n} + \ln n)^2) \quad \Theta(2^{n+\lg n}).$$

- (a) Which are equal?
- (b) Arrange the distinct classes in order from slowest growing to fastest growing. In other words, if  $\Theta(f(n))$  is to the left of  $\Theta(g(n))$ , then  $f(n) \in o(g(n))$ .

3. (30 pts.) The average running time for an algorithm is a nondecreasing function of  $n$  and satisfies  $T(4n) = T(2n) + 2T(n)$  for all  $n > 0$ . Furthermore,  $T(1) = 1$  and  $T(2) = 3$ .

(a) Determine  $T(2^k)$  as a function of the integer  $k$ .

*Hint:* Set  $t_k = T(2^k)$ .

(b) Determine the complexity class of  $T(n)$ .

4. (30 pts.) Suppose we have two sorted lists  $a_1, \dots, a_n$  and  $b_1, \dots, b_n$ , both of length  $n$ , that we want to merge to obtain a sorted list of length  $2n$ , say  $c_1, \dots, c_{2n}$ . To do this, we must decide where the  $a_i$ 's fit among the  $b_j$ 's to produce the  $c$  list. The number of choices for this is  $\binom{2n}{n} \geq 4^n / (2n^{1/2})$ .

Suppose the merge is done comparisons of keys. Using the above information, derive a lower bound for the worst case number of key comparisons that are needed. Explain your reasoning; don't just give an answer.

5. (30 pts.) Here is an informal description of a routine `Proc` that is looking for  $x$  in a sorted list  $S$ . The parameters are the ends of the list. While it is looking it does some processing in `ProcLow` and `ProcHigh`.

```

Proc(lo,hi)
  If lo > hi we are done.
  k = ⌊(lo + hi)/2⌋.
  If S[k] = x, we are done.
  If S[k] < x
    Call ProcHigh(k,hi) and Proc(k + 1,hi)
  Else
    Call ProcLow(lo,k) and Proc(lo,k - 1)
  Endif.
End

```

We begin by calling `Proc(1, n)`. Most of the time is spent in `ProcLow` and `ProcHigh`. In fact, `ProcLow(a, b)` requires  $\lg(b - a + 1)$  basic operations and `ProcHigh(a, b)` requires  $(b - a + 1)$  basic operations. (You do *not* need to know what any of this code is supposed to do.)

- (a) Let  $W(n)$  be the worst case running time for `Proc(1, n)`. Give a recursion and initial condition for  $W(2^n)$ . (In the worst case,  $x$  is not in the list.)
- (b) Let  $A(n)$  be the average running time for `Proc(1, n)`. Assuming  $x$  is not in the list and the probability that  $S[k] < x$  is  $1/2$ , give a recursion for  $A(n)$ . You need *not* give an initial condition.

6. (65 pts.) Indicate whether true or false. Beware of guessing:

correct answer +5pts.    incorrect answer -3pts.    no answer 0pts

\_\_\_  $\Theta(2^{n+2}) = \Theta(2^n)$ .

\_\_\_  $\Theta((n+2)^2) = \Theta(n^2)$ .

\_\_\_  $\Theta(2^{n+\lg n}) = \Theta(2^n)$ .

\_\_\_  $\Theta((n+\lg n)^2) = \Theta(n^2)$ .

\_\_\_ Greedy algorithms are called “greedy” because they make the best choice at the present time, without concern for the future.

\_\_\_ Dynamic programming algorithms use a bottom up approach.

\_\_\_ Divide and conquer algorithms use a bottom up approach.

\_\_\_ If a divide and conquer algorithm requires recomputing the same quantity many times, it is a good idea to look for a dynamic programming algorithm.

\_\_\_ No greedy algorithm is known for the 0-1 Knapsack Problem.

\_\_\_ It is usually fairly easy to determine average and worst-case time complexities for backtracking algorithms.

\_\_\_ There is a search algorithm that uses comparison of keys and is significantly faster on average and in the worst case than binary search.

\_\_\_ There is a sorting algorithm that uses comparison of keys and is significantly faster on average than and in the worst case than mergesort.

\_\_\_ Quicksort has a good average run time and a poor worst-case run time.