

1. (45 pts.) (a) Compute the rank of the permutation 1, 5, 3, 4, 2 when the permutations of $\{1, 2, 3, 4, 5\}$ are listed in direct insertion order.

Ans. If we start numbering positions with 0 and let $p(i)$ be the position where i is inserted, then the rank is $(5!/2!)p(2) + (5!/3!)p(3) + (5!/4!)p(4) + (5!/5!)p(5)$. Since $p(2) = 0$, $p(3) = 1$, $p(4) = 1$, and $p(5) = 3$, the rank is 28. You could also draw the relevant part of the decision tree here and in (b) and (c) to figure things out.

- (b) What permutation immediately follows 1, 5, 3, 4, 2 in direct insertion order?

Ans. Increase $p(5)$ by 1, giving 5, 1, 3, 4, 2.

- (c) What permutation immediately follows 1, 5, 3, 4, 2 in lex order?

Ans. Starting from the right, we look for the first position that can be increased. The 2 can't be changed since there's nothing to replace it with. The 4 can only be replaced by 2, which is a decrease. The 3 can be replaced by 2 or 4. To increase, use 4. Next arrange the left over 2 and 3 to give the leftmost possibility. We have 1, 5, 4, 2, 3.

2. (20 pts.) Find the unlabeled rooted plane tree with 7 leaves and rank 60.

Ans. We have $56 = b_1b_6 + b_2b_5 \leq 60 < b_1b_6 + b_2b_5 + b_3b_4$. Also $60 - 56 = 0b_4 + 4$. Thus the left subtree of the root has 3 leaves and rank 0 while the right has 4 leaves and rank 4. Each of these is easy to determine: Rank 0 means keep all the branching to the right. Rank $b_k - 1$ means keep all the branching to the left.

$$b_1 = b_2 = 1, \quad b_3 = 2, \quad b_4 = 5, \quad b_5 = 14, \quad b_6 = 42, \quad b_7 = 132.$$

3. (30 pts.) Let $C_n(k)$ be the number of times position k is changed in the Gray code for n -long vectors of zeroes and ones given in the text. You may use the following fact without proving it:

$$C_n(k) = \begin{cases} 1 & \text{if } k = 1, \\ 2C_{n-1}(k-1) & \text{if } 1 < k \leq n, \\ 0 & \text{if } k > n. \end{cases}$$

- (a) Tabulate values of $C_n(k)$ for $1 \leq k \leq n \leq 4$.

Ans. This is straightforward.

- (b) State and prove a simple formula for $C_n(k)$ that does not involve a recursion.

Ans. You can guess from (a) that $C_n(k) = 2^{k-1}$ for $1 \leq k \leq n$ and zero otherwise. This can be proved by induction. Should we induct on n or k ? Either will work. However, since the answer does not depend on n and " $C_n(k) = 1$ if $k = 1$ " looks like an initial condition, it may be somewhat easier to induct on k . (It is.) As noted, this is the base case. For k in the range $2 \leq k \leq n$, we have

$$\begin{aligned} C_n(k) &= 2C_{n-1}(k-1) && \text{by the given formula,} \\ &= 2 \times 2^{k-2} && \text{by } \mathcal{A}(k-1), \\ &= 2^{k-1}. \end{aligned}$$

4. (30 pts.) Let b_n be the number of binary unlabeled rooted plane trees with n leaves. It is known that $b_n < 4b_{n-1}$ for $n > 1$ and you may use this fact without proof. (It can be proved by using Exercise 9.1.12.)

- (a) Show that, for more than $b_n/4$ of these trees with n leaves, the left subtree consists of just a single leaf when $n > 1$.

In other words, show that if the two edges leading from the root go to trees T_1 and T_2 , then we have $|T_1| = 1$ in more than $b_n/4$ of the cases.

Hint: What does the term $b_i b_{n-i}$ in $b_n = b_1 b_{n-1} + b_2 b_{n-2} + \cdots + b_{n-1} b_1$ count?

Ans. Let $|T|$ be the number of leaves of T . Following the hint, $b_i b_{n-i}$ counts those T with $|T| = n$, $|T_1| = i$ and $|T_2| = n - i$. Thus the number with $|T_1| = 1$ is $b_1 b_{n-1} = b_{n-1} > b_n/4$, where the inequality follows from $B_n < 4b_{n-1}$.

- (b) Find a constant $P > 1/4$ such that the following statement is true for $n > 1$ about those binary trees with n leaves.

“In more than Pb_n of them, the left subtree contains at most 2 leaves.”

Ans. From $b_N < 4b_{N-1}$ at $N = n$ and $N = n - 1$, we have $b_n < 4b_{n-1} < 4^2 b_{n-2}$. Thus $b_{n-2} > b_n/16$. Reasoning as in (a), the number of such trees is greater than

$$b_1 b_{n-1} + b_2 b_{n-2} = b_{n-1} + b_{n-2} < b_n/4 + b_n/16 = 5b_n/16.$$

Thus P can have any value such that $1/4 < P \leq 5/16$.