

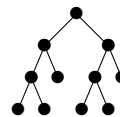
1. Let $G = (V, E)$ where $V = \{0, 1, a, b, A, B\}$ and $E = \{\{0, 1\}, \{0, a\}, \{0, b\}, \{0, A\}, \{0, B\}, \{a, b\}, \{A, B\}\}$.


Sketch the simple graph G and **compute** its chromatic polynomial.

- A.** We omit the sketch. There are various ways to compute the chromatic polynomial. One can use formulas, but the easiest is to appeal directly to the definition. Given x colors, there are x ways to color vertex 0. There are then $x - 1$ ways to color each of 1, a , and A since they must differ from 0. There are then $x - 2$ ways to color each of b and B . (For example, b must differ from 0 and a , which leaves $x - 2$ colors for it.) By the Rule of Product from Chapter 1, $P_G(x) = x(x - 1)^3(x - 2)^2$.

2. **Compute** the rank of the binary RP-tree shown here.

For your information, $b_1 = b_2 = 1$, $b_3 = 2$, $b_4 = 5$, $b_5 = 14$, $b_6 = 42$, and $b_7 = 132$.

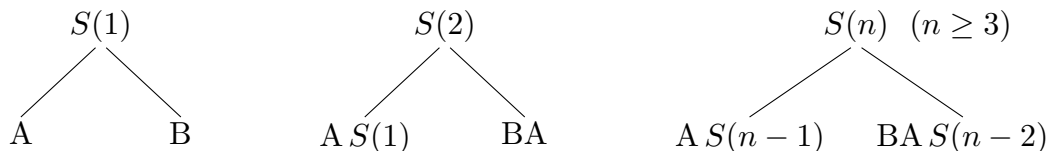


- A.** Let B be the given tree and let $T =$ . Then

$$\text{RANK}(B) = b_1 b_5 + b_2 b_4 + \text{RANK}(T) b_3 + \text{RANK}(T).$$

You can compute $\text{RANK}(T)$ or note that T is the second of the two trees counted by b_3 and so has rank 1. Thus the answer is $14 + 5 + 2 + 1 = 22$.

3. The local description of a decision tree for constructing sequences of A's and B's is given below. The notation $BA S(n - 2)$ means place BA in front of each sequence produced by $S(n - 2)$.



Let $S^*(n)$ denote the entire decision tree. Thus $S^*(1) = S(1)$ and $S^*(2)$ has the three leaves AA, AB, and BA.

- (a) **Find** a recursion for s_n , the number of leaves of $S^*(n)$.

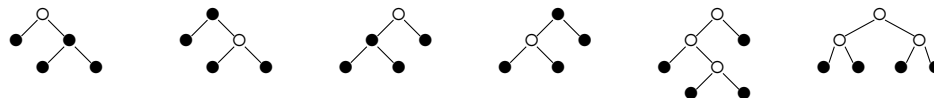
Remember to include initial conditions.

- A.** From the pictures, we read off $s_1 = 2$, $s_2 = s_1 + 1 (= 3, \text{ if you wish})$, and $s_n = s_{n-1} + s_{n-2}$ for $n \geq 3$.

(b) **Prove** that the leaves of $S^*(n)$ are sequences of length n and that their order from left to right is alphabetic.

A. We prove it by induction on n . For $n = 1$, it is clear from the picture. For $n = 2$, the entire tree has leaves AA, AB, and BA from left to right. Now we use induction for $n \geq 3$. Since A precedes B, the leaves of $A S^*(n-1)$ precede the leaves of $BA S^*(n-2)$ in alphabetic order. The leaves of $S^*(n-1)$ all have length $n-1$ and are in alphabetic order. Thus the leaves of $A S^*(n-1)$ all have length n and are in alphabetic order. The leaves of $S^*(n-2)$ all have length $n-2$ and are in alphabetic order. Thus the leaves of $BA S^*(n-2)$ all have length n and are in alphabetic order. This completes the proof.

4. A binary RP-tree has information stored at each leaf vertex. Each non-leaf vertex may or may not have information stored at it. Let t_n be the number of such trees with information stored at exactly n vertices and let $T(x) = \sum t_n x^n$ be the generating function. The following picture shows some of the nine trees that contribute to t_4 . An empty circle indicates a vertex with no information.



Find a formula for $T(x)$ similar to the formula $B(x) = x + B(x)^2$ we found for binary RP-trees. *To receive credit you must justify your formula; that is, explain how you got it.*

- A.** A tree of the desired type is either
- (a) a single vertex with information (since it is a leaf and so has information) OR
 - (b) a root with information joined to two trees of the same type OR
 - (c) a root without information joined to two trees of the same type.

The generating function for a vertex with information is $x^1 = x$ and that for a vertex without information is $x^0 = 1$. Using the Rule of Product in (b) and (c) and the Rule of Sum to combine the results we have

$$T(x) = x + xT(x)^2 + 1T(x)^2 = x + (x + 1)T(x)^2.$$