# Density (Ramsey) theorems for intersection graphs of $t$-monotone curves 

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## We will only consider simple topological graphs.



## Three conjectures in topological graph theory.

conjecture 1: Thrackle conjecture.

## Conjecture (Conway)

Every n-vertex simple topological graph with no two disjoint edges, has at most $n$ edges.

Fulek and Pach 2010: $|E(G)| \leq 1.43 n$.

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If edges are segments: Yes, Erdős.
If edges are x-monotone: Yes, Pach and Sterling 2011.

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conjecture 2: Extremal problem (generalization):

## Conjecture (Pach and Tóth 2005, sparse graphs)

Every n-vertex simple topological graph with no $k$ pairwise disjoint edges has at most $c_{k} n$ edges.

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If edges are segments: Yes, Pach and Töröcsik 1993.
If edges are x-monotone: Yes, Pach and Töröcsik 1993.
Note: Every complete $n$-vertex simple topological graph has $\Omega\left(n^{1 / 3}\right)$ pairwise disjoint edges, Suk 2011.

All solved for $x$-monotone curves, but all are still open for 2-monotone curves.

## Conjecture (Trackle)

Every n-vertex simple topological graph with no two disjoint edges, has at most $n$ edges.

## Conjecture (Sparse problem)

Every n-vertex simple topological graph with no $k$ pairwise disjoint edges has at most $c_{k} n$ edges.

## Conjecture (dense problem)

Every $n$-vertex simple topological graph with $\Omega\left(n^{2}\right)$ edges, has $n^{\delta}$ pairwise disjoint edges.

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A curve $\gamma$ is $t$-monotone if its interior has at most $t-1$ vertical tangent points. 1 -monotone $=x$-monotone.

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## Results

Pach and Tóth's problem:

## Theorem (Suk 2012)

Let $G$ be an n-vertex simple topological graph with edges drawn as $t$-monotone curves. If $G$ has no $k$ pairwise disjoint edges, then $|E(G)| \leq n(\log n)^{c_{t} \log k}$.

Recall Pach and Tóth's bound of $n(\log n)^{4 k-8}$ for general curves.

## Results

## Corollary (Suk 2012)

Let $G$ be an n-vertex simple topological graph with edges drawn as $t$-monotone curves. If $|E(G)| \geq \Omega\left(n^{2}\right)$, then $G$ contains $n^{\delta_{t} / \log \log n}$ pairwise disjoint edges.

Fox and Sudakov showed $\log ^{1.02} n$ pairwise disjoint edges in the general case.

## Conjecture

Every n-vertex simple topological graph with at least $\Omega\left(n^{2}\right)$ edges, has $n^{\delta}$ pairwise disjoint edges.

## Theorem (Suk 2012)

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## Ramsey type result

## Theorem (Two color, Suk 2012)

Let $R$ be a family of $n$ red $t$-monotone curves in the plane, and let $B$ be a family of $n$ blue $t$-monotone curves in the plane, such that $R \cup B$ is simple. Then there exist subfamilies $R^{\prime} \subset R$ and $B^{\prime} \subset B$ such that $\left|R^{\prime}\right|,\left|B^{\prime}\right| \geq \epsilon n$, and either
(1) every red curve in $R^{\prime}$ intersects every blue curve in $B^{\prime}$, or
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(2) every red curve in $R^{\prime}$ is disjoint to every blue curve in $B^{\prime}$.
(1) For segments, Pach and Solymosi 2001.
(2) Semi-algebraic sets in $\mathbb{R}^{d}$, Alon et al. 2005.
(3) Definable sets belonging to some fixed definable family of sets in an o-minimal structure, Basu 2010.

All previous results assumed some type of bounded/fixed complexity.

Two color theorem + Szemerédi's regularity lemma $\Rightarrow$ density theorem $\Rightarrow$

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## Proof:



Select a random sample of $c$ blue curves, for large constant $c$.


Trapezoid decomposition of $\mathbb{R}^{2}$ : Draw a vertical line through each endpoint and through each vertical tangent point.


At most $c_{t}^{2}$ number of cells. With high probability, each cell will intersect at most $n / 2$ blue curves!

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By pigeonhole, there exists a cell with at least $\epsilon n$ number of "left-endpoints", and $n / 2$ blue curves is disjoint to this cell.




Look at the remaining red and blue curves.


We have $\epsilon n$ red curves, and $\epsilon n$ blue curves remaining.

Do it again for the remaining right-endpoints and the remaining blue curves.


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In the end, we have $\delta n$ red curves, and two regions $R_{1}$ and $R_{2}$ that contains the endpoint of these red curves.


And no blue curves intersects the interior of $R_{1}$ and $R_{2}$.


Do this whole process again with the endpoints of the blue curves to get regions $R_{3}$ and $R_{4}$.


Endpoints of the blue curves lie inside $R_{3}$ and $R_{4}$, and all red curves are disjoint to the interior of $R_{3}$ and $R_{4}$.


Apply a case analysis/Jordan curve argument to find:


End of "proof".

## Open problem

$t$-monotone condition only used for the trapezoid decomposition.

## Problem

Given an n-point set $P$ and family $F$ of $n$ simple curves, such that no point lies on any curve in $F$, does there exist a region $R$ that contain $\epsilon$ points, and the interior of $R$ is disjoint to $\epsilon n$ curves from $F$ ?


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## Problem (2-monotone thrackle)

Let $G$ be an n-vertex simple topological graph with edges drawn as 2-monotone curves. If $G$ does not contain 2 disjoint edges, then $|E(G)| \leq n$ ?

## Problem (2-monotone color)

Given a simple family $F$ of 2-monotone curves in the plane with no 3 pairwise disjoint members, $\chi(\bar{F}) \leq c$ for some constant $c$ ?

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Note: Color problem is true for segments/ $x$-monotone curves.

## Thank you!

