

Lecture 2: Russell's example.

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Ex. This is a false sentence.

• It is NOT a proposition.

Suppose to the contrary that it is a proposition.

(This method of proof is called proof by contradiction.)

Then there are two possibilities:

(This method of arguing is called case-by-case proof.)

<p><u>Case 1.</u> It is a <u>true</u> proposition.</p> <p>So the claim of this proposition is supposed to be true, which says it is false. That is a contradiction.</p>	<p><u>Case 2.</u> It is a <u>false</u> proposition.</p> <p>So the claim of this proposition is supposed to be false, which implies that it is true. That is a contradiction.</p>
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The only reason that we are getting a contradiction is because we assumed that the above sentence is a proposition. Hence this is NOT a valid assumption, i.e. the above sentence is NOT a proposition.

Lecture 2: Propositional forms and truth tables

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A propositional form is a legitimate expression involving logical variables, connectives $\wedge, \vee, \Rightarrow, \Leftrightarrow, \neg$, and $(,)$.

Here are the truth table of conjunction \wedge , disjunction \vee , and implication \Rightarrow .

P and Q , $P \wedge Q$

all the possibilities of the truth values of P, Q .

P	Q	$P \wedge Q$
T	T	T
T	F	F
F	T	F
F	F	F

It is "and" so both should hold for P and Q to hold

If there are 3 variables, the number of rows is 8. For 4 variables, we get 16 rows. For n variables, we get 2^n rows.

P or Q , $P \vee Q$

This is slightly different from the way "or" is used in a daily language. There we sometimes assume P and Q do not hold at the same time.

P	Q	$P \vee Q$
T	T	T
T	F	T
F	T	T
F	F	F

For this to fail, both P and Q should fail

P implies Q , $P \Rightarrow Q$

Implication fails only when hypothesis (P) holds and conclusion (Q) fails.

P	Q	$P \Rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

In particular $P \Rightarrow Q$ is true if P is false.

Lecture 2: Equivalence of propositional forms

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Ex. $(1=2) \Rightarrow$ Sun is a moon.

is a true proposition as the hypothesis is false.

Of course I am NOT claiming that the conclusion "sun is a moon"

is correct. We are only saying that the implication is true.

Ex. Write the truth table of $\neg(P \vee Q)$ and $(\neg P) \wedge (\neg Q)$.

P	Q	$P \vee Q$	$\neg(P \vee Q)$	$\neg P$	$\neg Q$	$(\neg P) \wedge (\neg Q)$
T	T	T	F	F	F	F
T	F	T	F	F	T	F
F	T	T	F	T	F	F
F	F	F	T	T	T	T

Definition Two propositional forms are called equivalent if they have the same truth tables.

Ex. $\neg(P \vee Q) \equiv (\neg P) \wedge (\neg Q)$ } de Morgan's law.

Similarly $\neg(P \wedge Q) \equiv (\neg P) \vee (\neg Q)$ }

Convention. T: a true proposition.

⊥: a false proposition. (contradiction).

Lecture 2: Proof by contradiction

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Ex. (Proof by contradiction) To prove P it is equivalent to assume $(\neg P)$ and reach to a contradiction.

$$P \equiv (\neg P) \Rightarrow \perp$$

P	$\neg P$	\perp	$(\neg P) \Rightarrow \perp$
T	F	F	T
F	T	F	F

(as the hypoth. is false)
(as the hyp. holds and the conclusion fails.)